

Network Flow

Network Flow



Maximal Flow Problem (Ford Fulkerson Rule)

1. The Ford Fulkerson method is used for solving maximum flow problem.

2. Basic terms:

- Source
- Sink
- Capacity or bottle neck capacity
- Flow
- Augmenting path
- Residual capacity

Source : The source vertex has all outward edge, no inward edge.

Sink : Sink will have all inward edge, no outward edge.

Bottleneck Capacity : Bottleneck Capacity of the path is the minimum capacity of any edge on the path.

Residual Capacity : Every edge of a residual graph has a value called residual Capacity, which is equal to original Capacity minus current flow.

Network flows

- A Network is represented by a weighted graph in which

vertices — stations

edges — lines

through which the given weight assigned
to each edge shows the capacity.

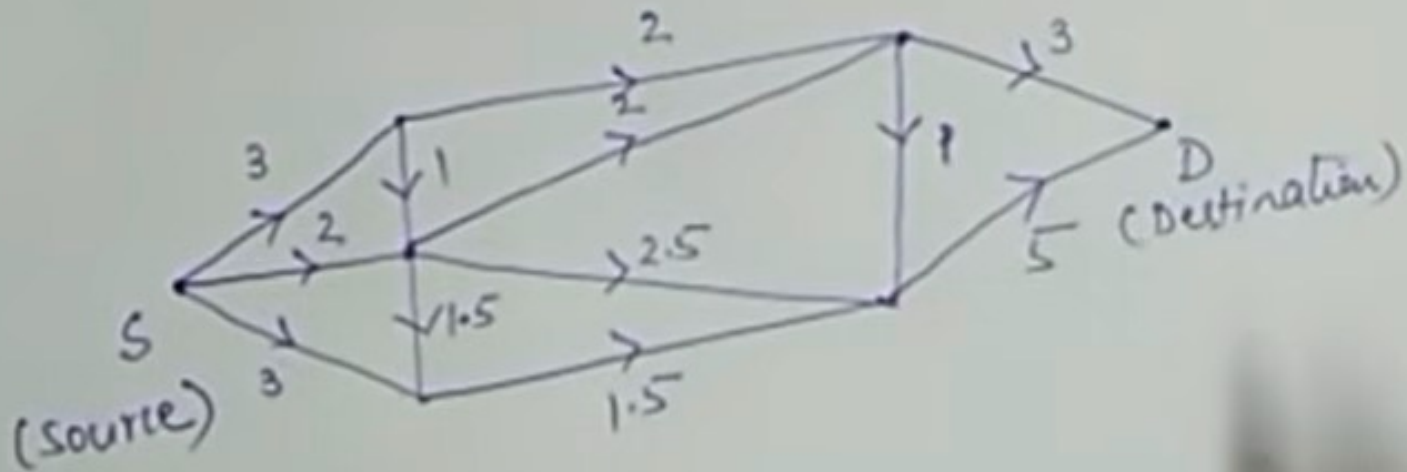
edges — lines

through which the given weight assigned
to each edge shows the Capacity.

- Capacity is the max amount of flow possible per unit of time through that line.

- Capacity is the max possible per unit of time through that line.

- A network has satisfies following properties:-
 - Contain one & only one vertex - no incoming edge.
 - Contain one & only one vertex - no outcoming edge.



Flow

$$a) \quad 0 \leq f(v_i, v_j) \leq c(v_i, v_j)$$

$$b) \quad \sum_{v_i} f(v_i, v_j) = \sum_{v_k} f(v_j, v_k)$$

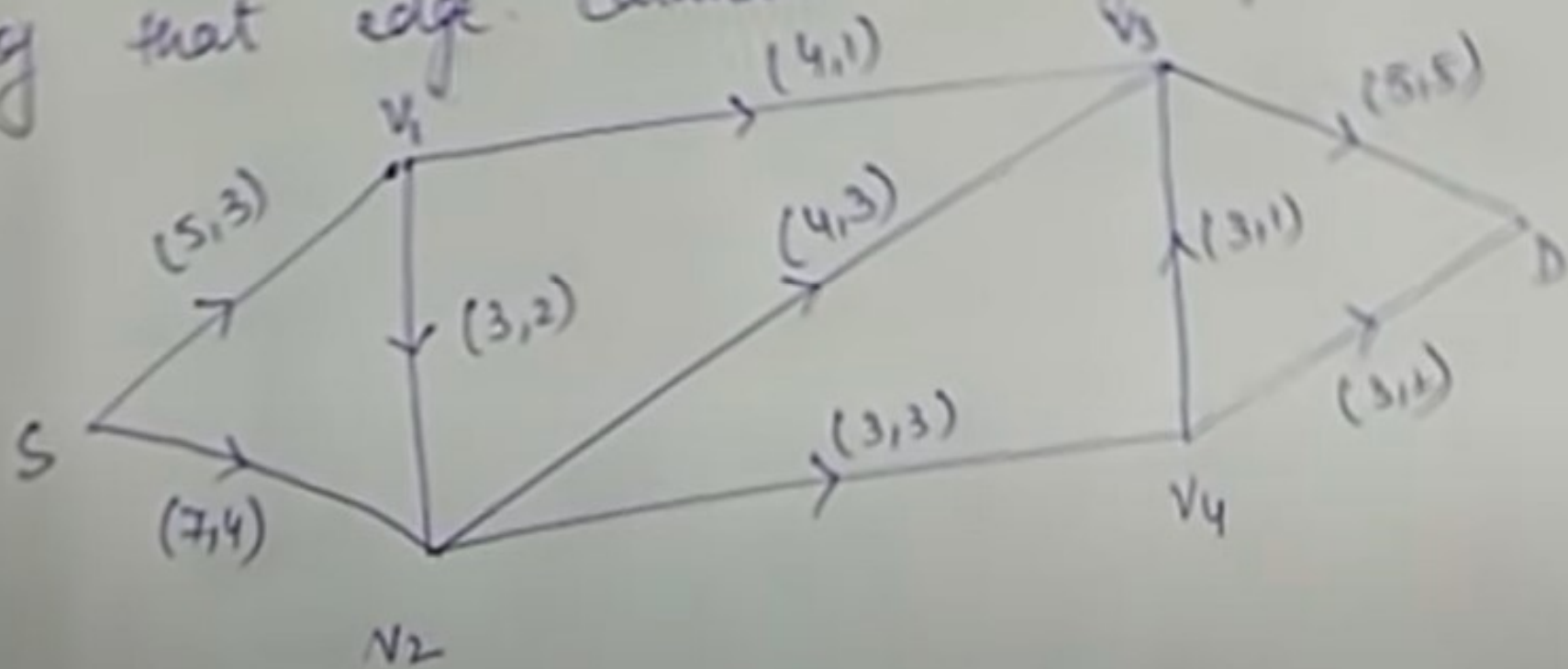
$$c) \quad f(v_i, v_j) = 0$$

Value of the flow :-

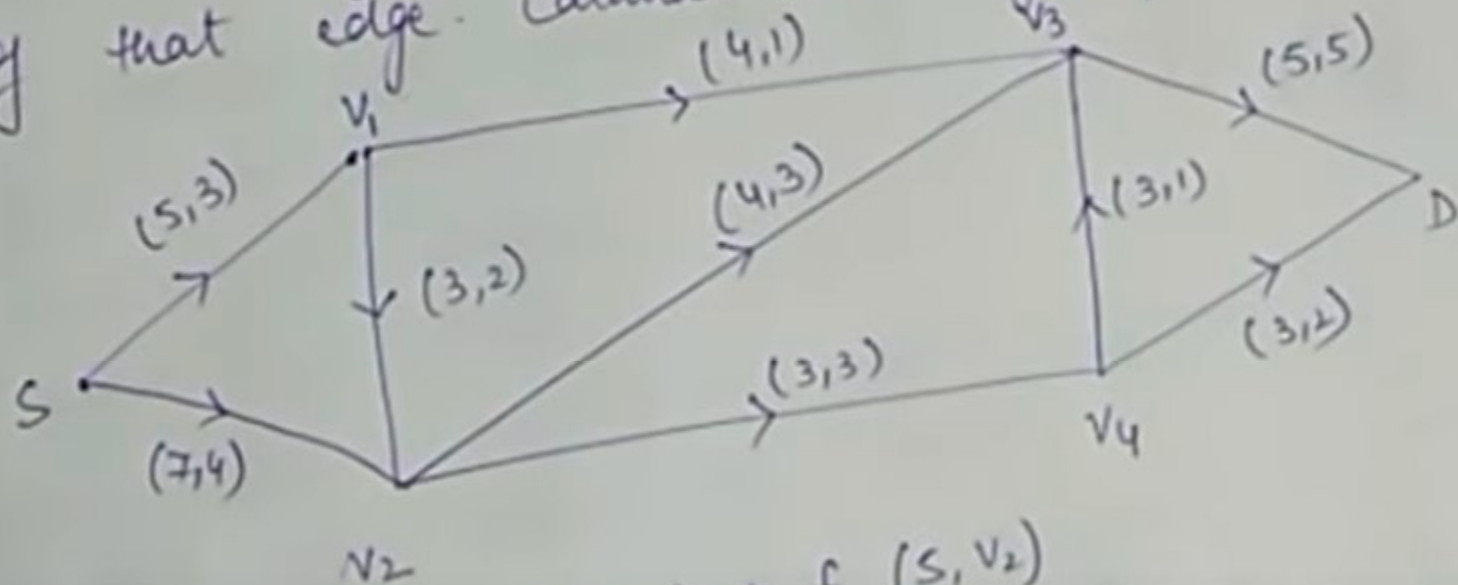
$$f_v = \left[\sum_{v_i} f(s, v_i) - \sum_{v_i} f(v_i, s) \right]$$
$$= \left[\sum_{v_k} f(v_k, D) - \sum_{v_k} f(D, v_k) \right]$$

Algorithms

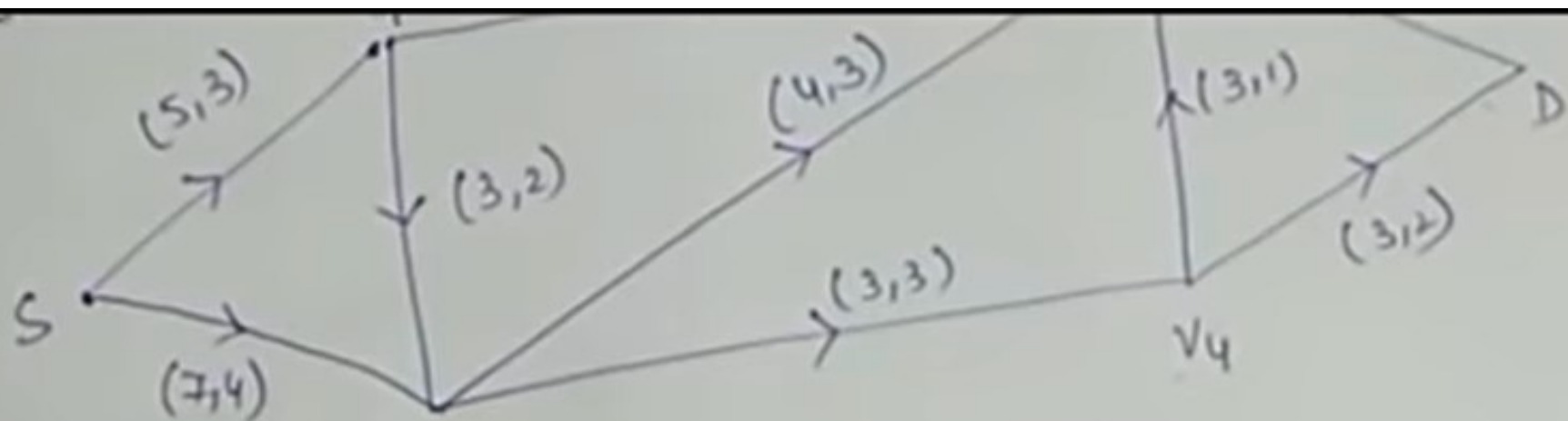
Consider a transport n/w. The pair (C, f) assigned to each edge indicate the capacity and flow of that edge. Calculate value of flow.



Q. Consider a transport network. To each edge indicate the capacity and flow of that edge. Calculate value of flow.



$$\begin{aligned}
 f_v &= f(S, v_1) + f(S, v_2) \\
 &= 3 + 4 \\
 &= 7
 \end{aligned}$$

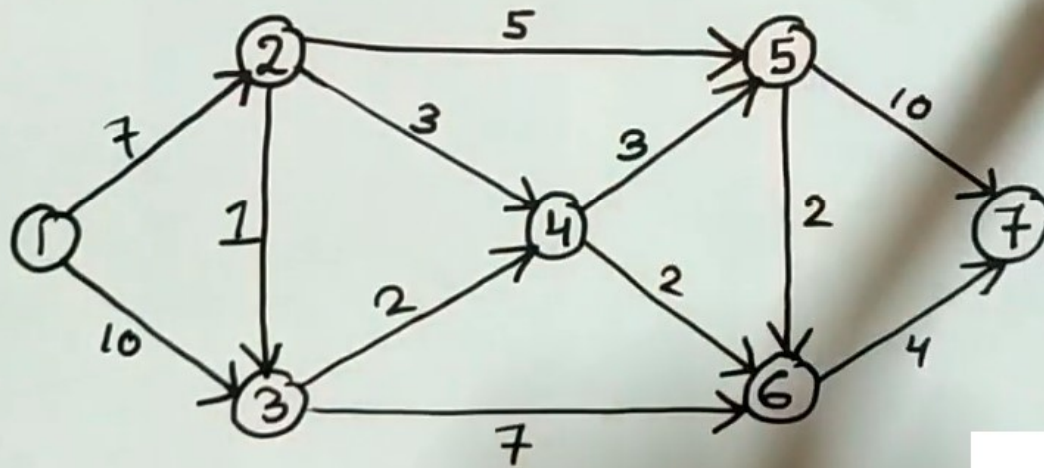


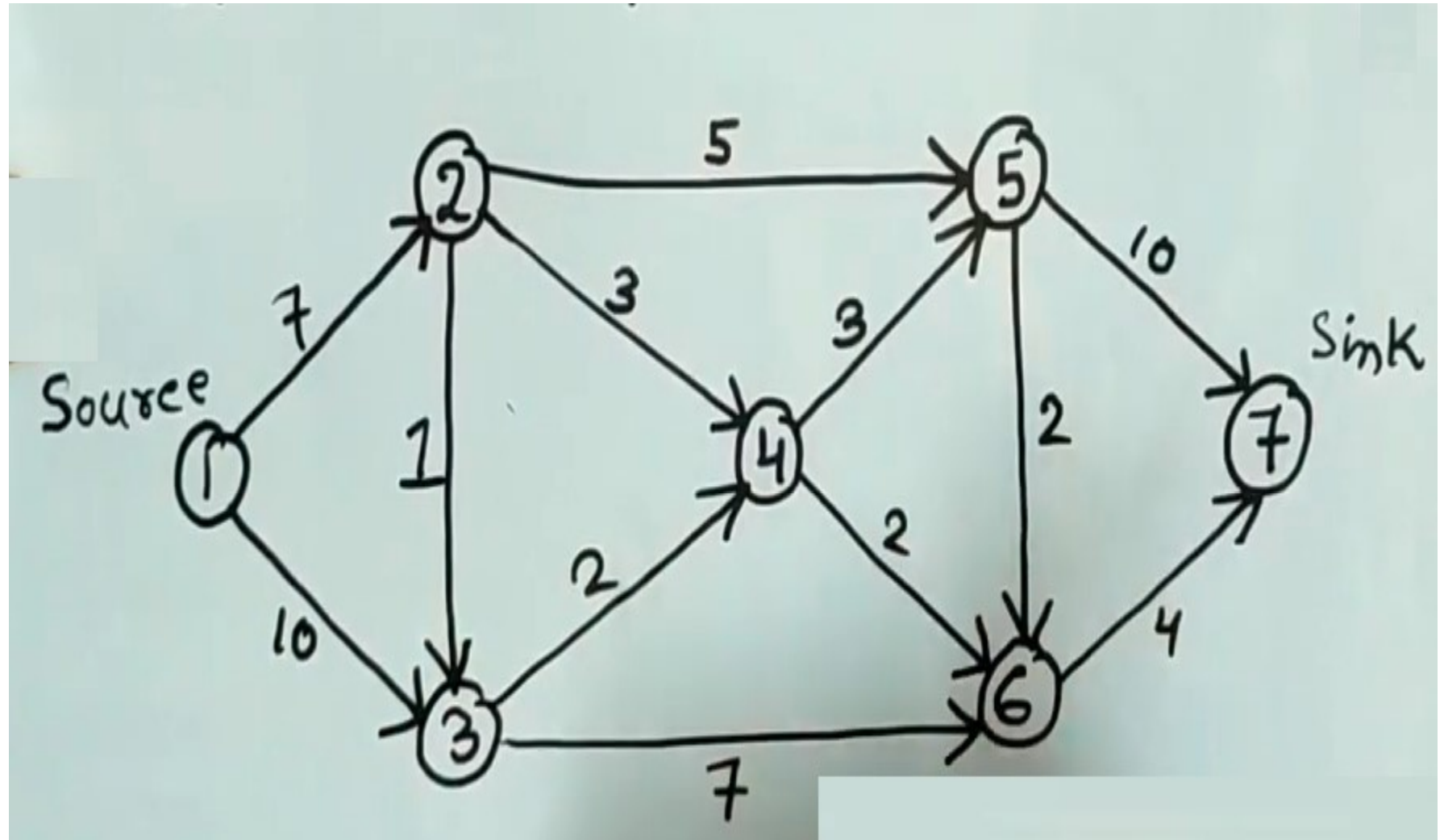
$$\begin{aligned}
 f_v &= f(S, v_1) + f(S, v_2) \\
 &= 3 + 4 \\
 &= \textcircled{7}
 \end{aligned}$$

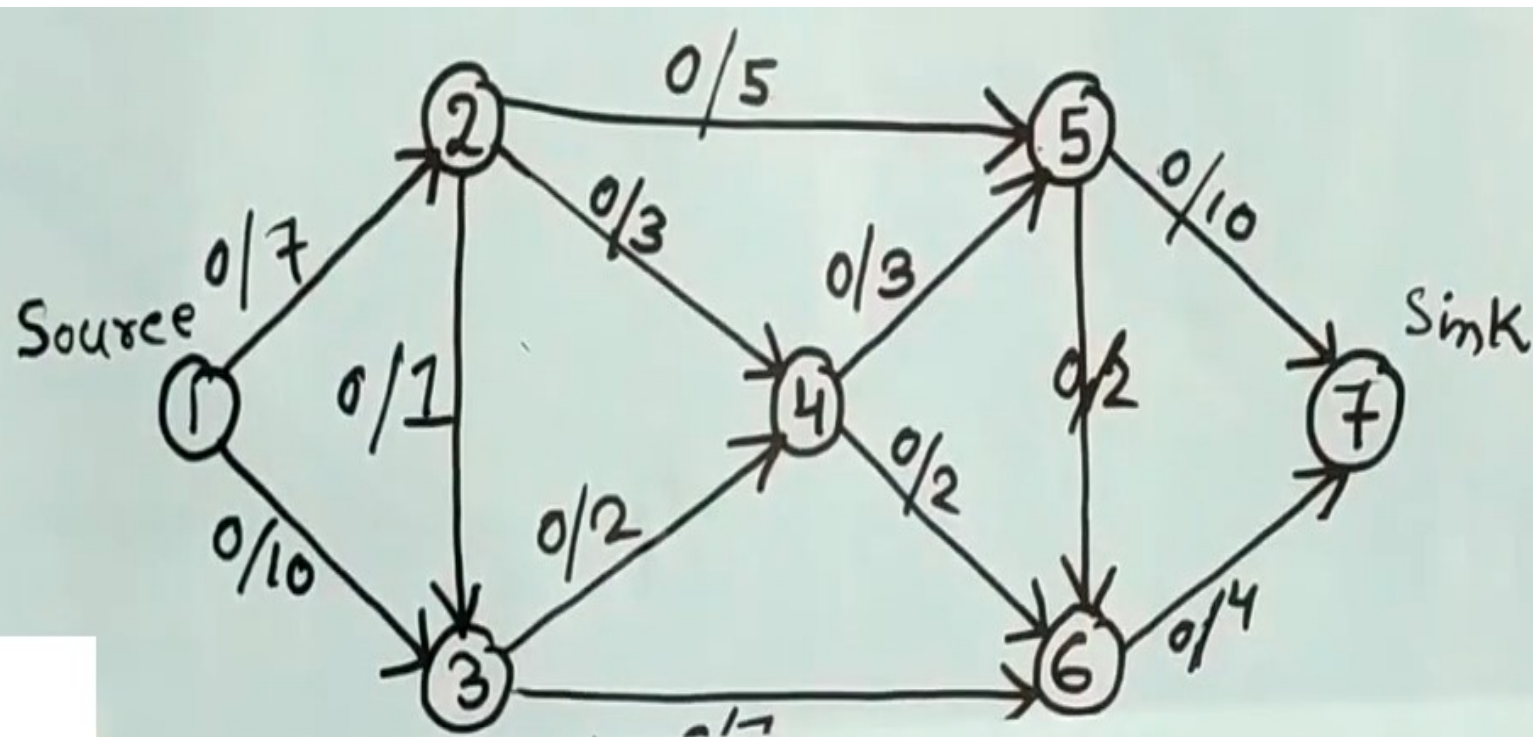
$$\begin{aligned}
 f_v &= f(v_3, D) + f(v_4, D) \\
 &= 5 + 2 \\
 &= \textcircled{7}
 \end{aligned}$$

Maximal Flow Problem (Ford Fulkerson Rule)

Q: Find the maximum flow through the given network using Ford Fulkerson algorithm.

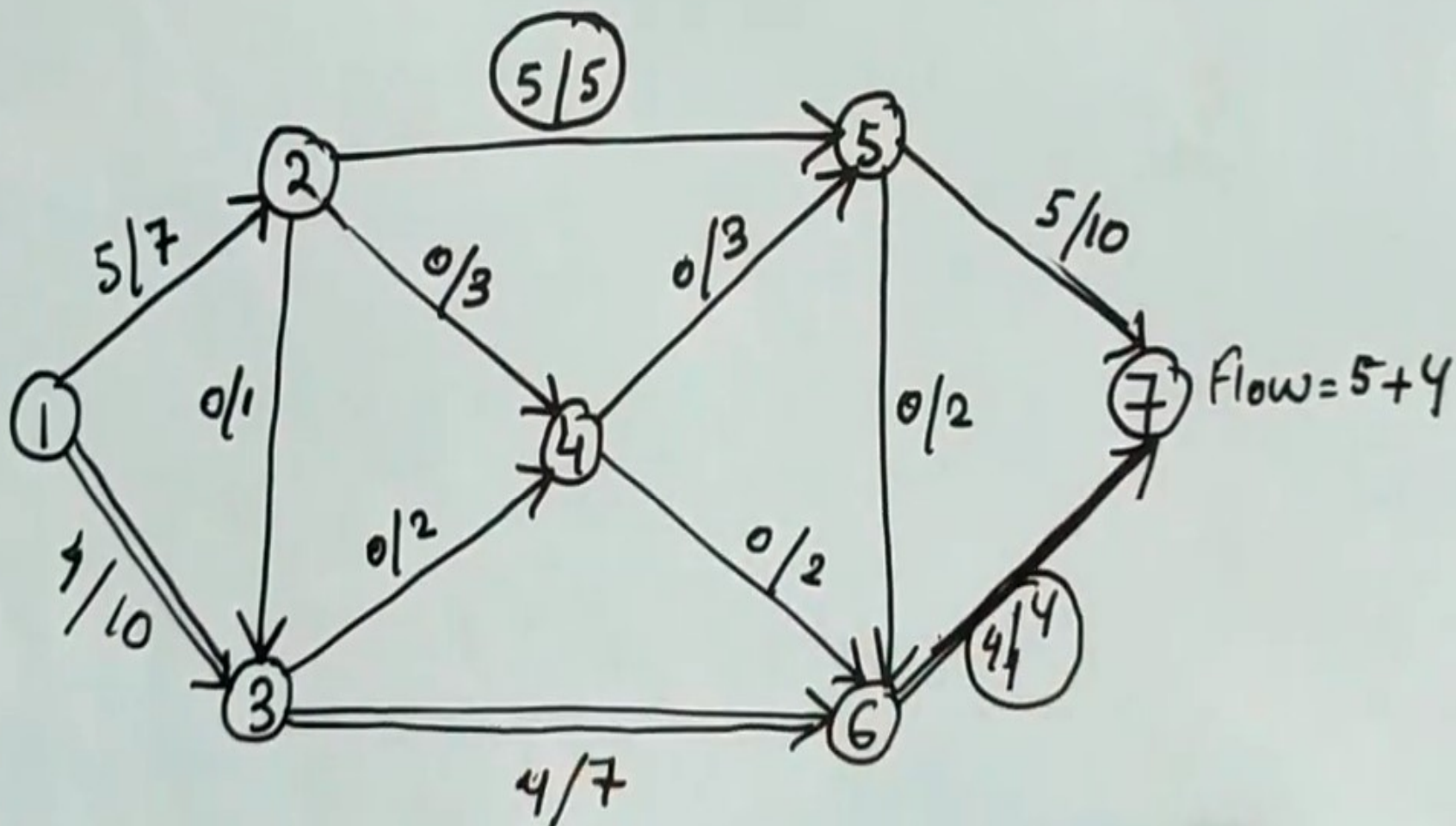


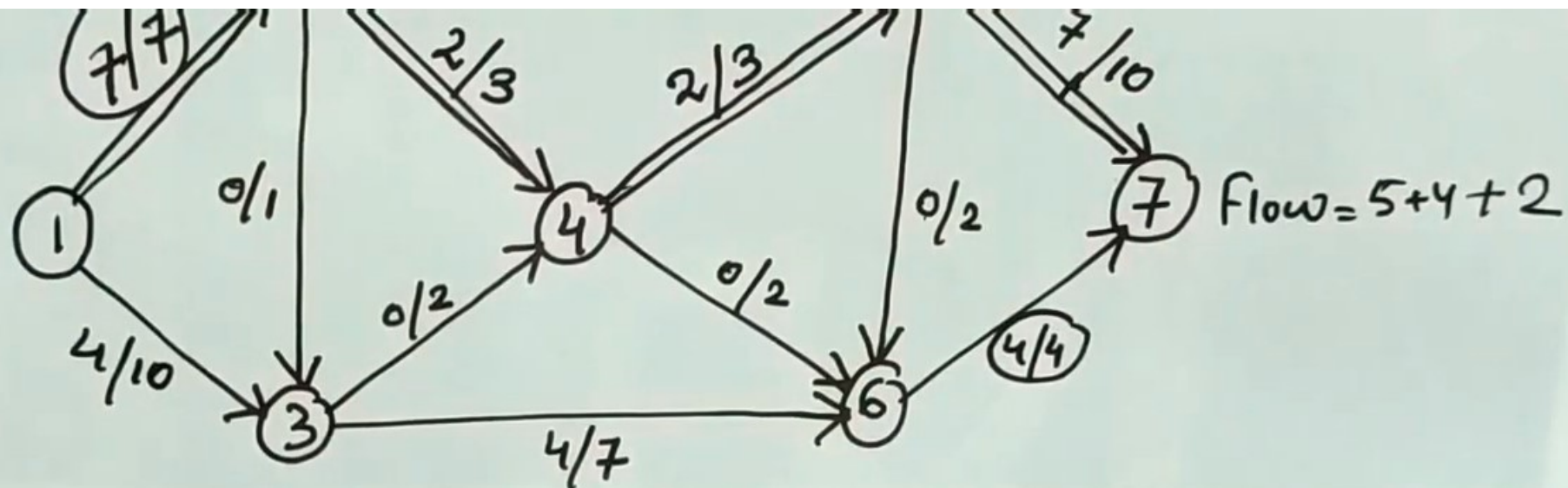




Flow
→ 4/5 ✓

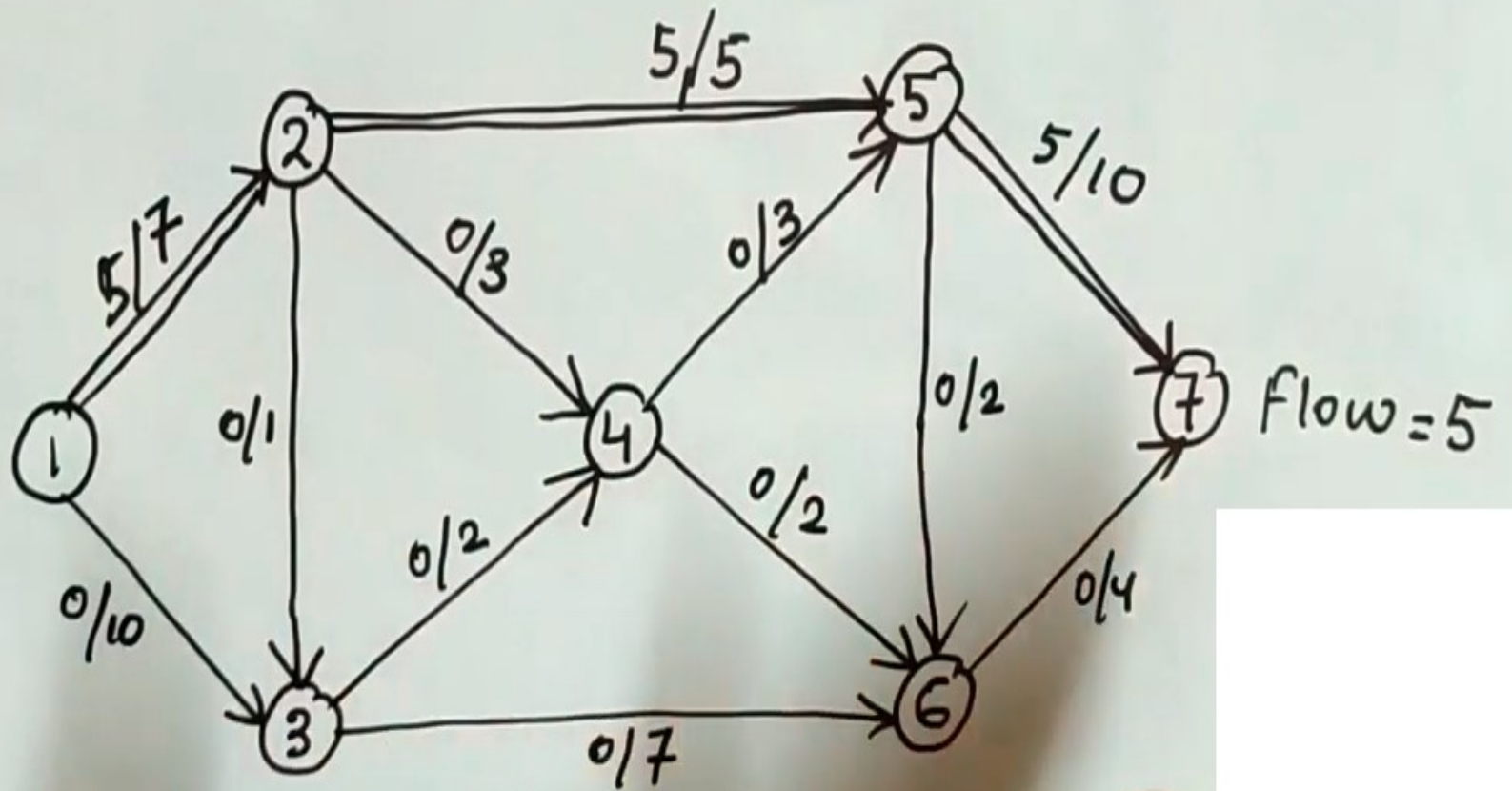
Augmenting path	Bottleneck capacity
1 - 2 - 5 - 7	5





Augmenting path	Bottleneck capacity
1 - 2 - 5 - 7	5
1 - 3 - 6 - 7	4
1 - 2 - 4 - 5 - 7	2

Solution :



Ford-Fulkerson Algorithm for Maximum Flow Problem

Problem

Given a graph which represents a flow network where every edge has a capacity. Also given two vertices Source s and sink t in the graph. Find out the maximum possible flow from s to t with following constraints.

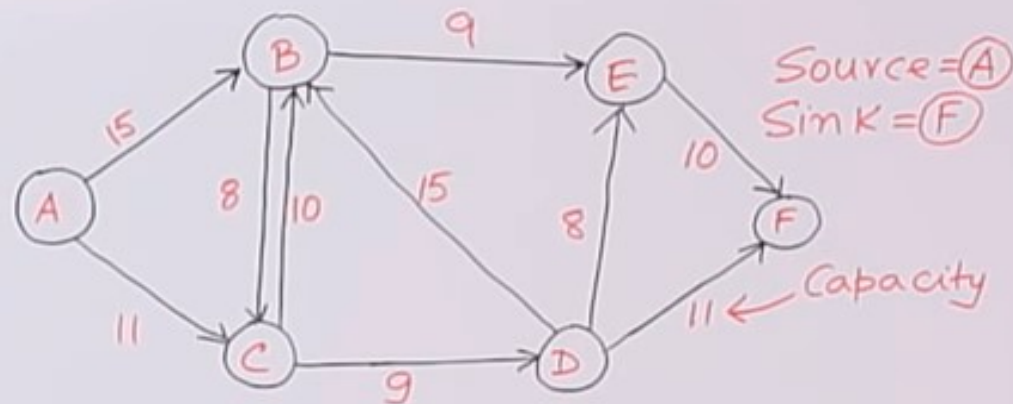
- a) Flow on an edge doesn't exceed the given capacity of the edge.
- b) In-flow is equal to Out-flow for every vertex except s and t .

Algorithm

Ford-Fulkerson Algorithm

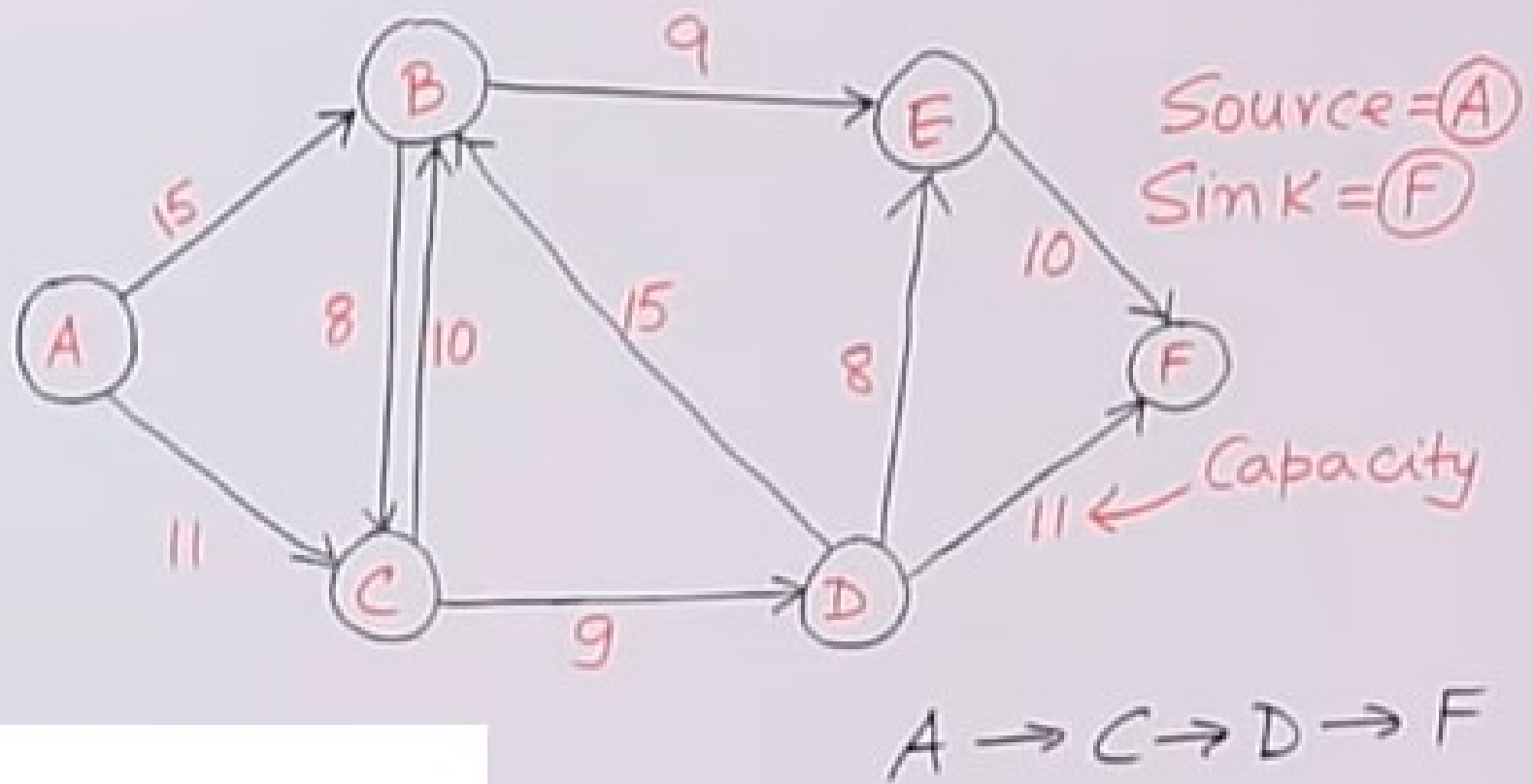
The following is a simple idea of the algorithm

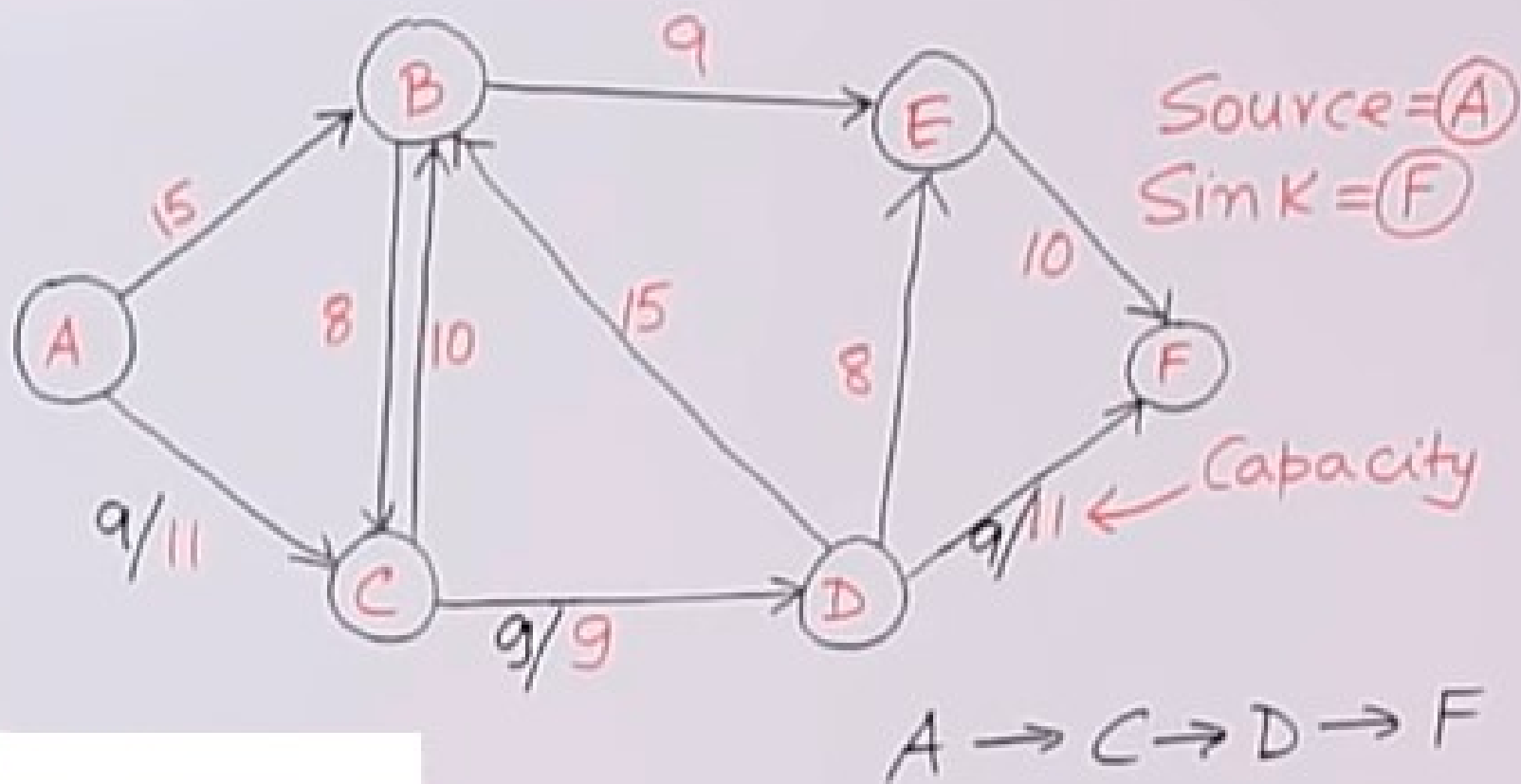
- 1) Start with a initial flow as 0.
- 2) While there is an augmenting path from source to sink
Add this path flow to flow
- 3) Return flow

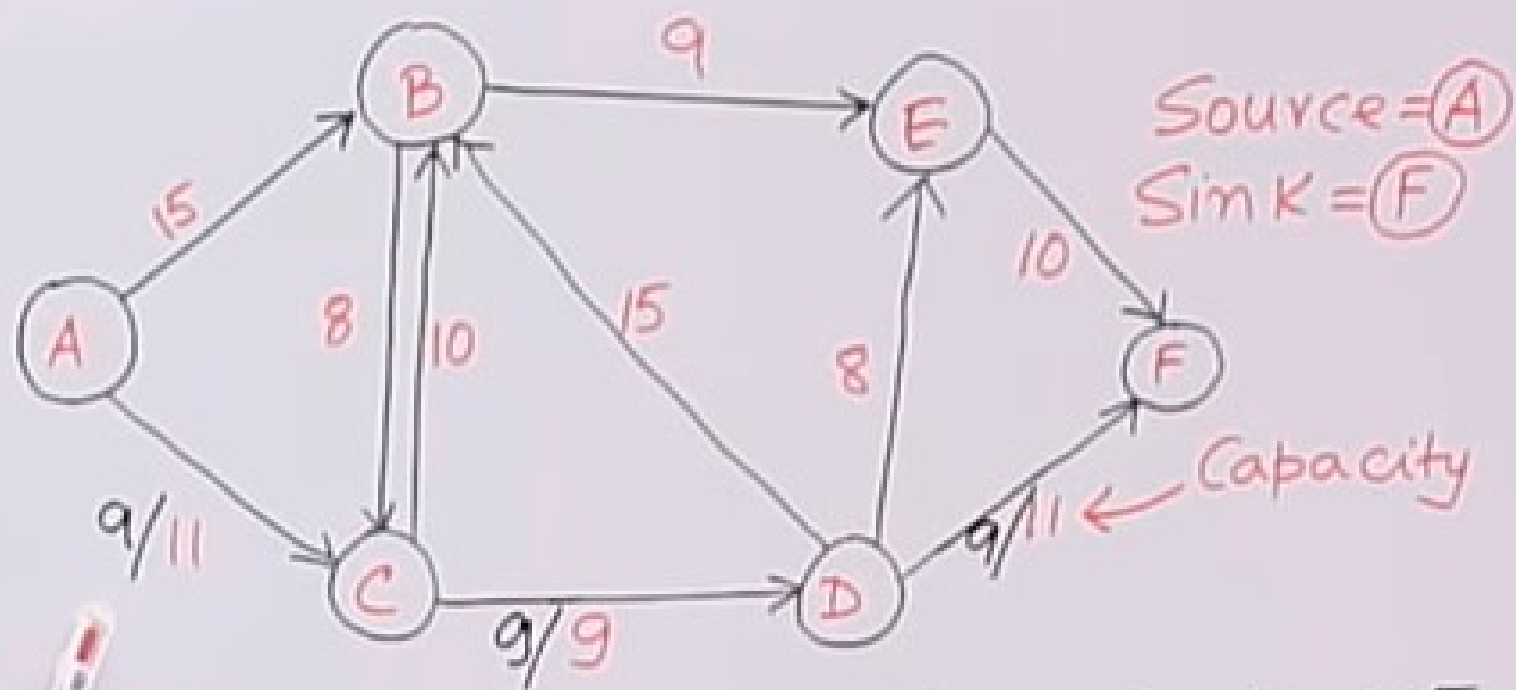


Terminologies

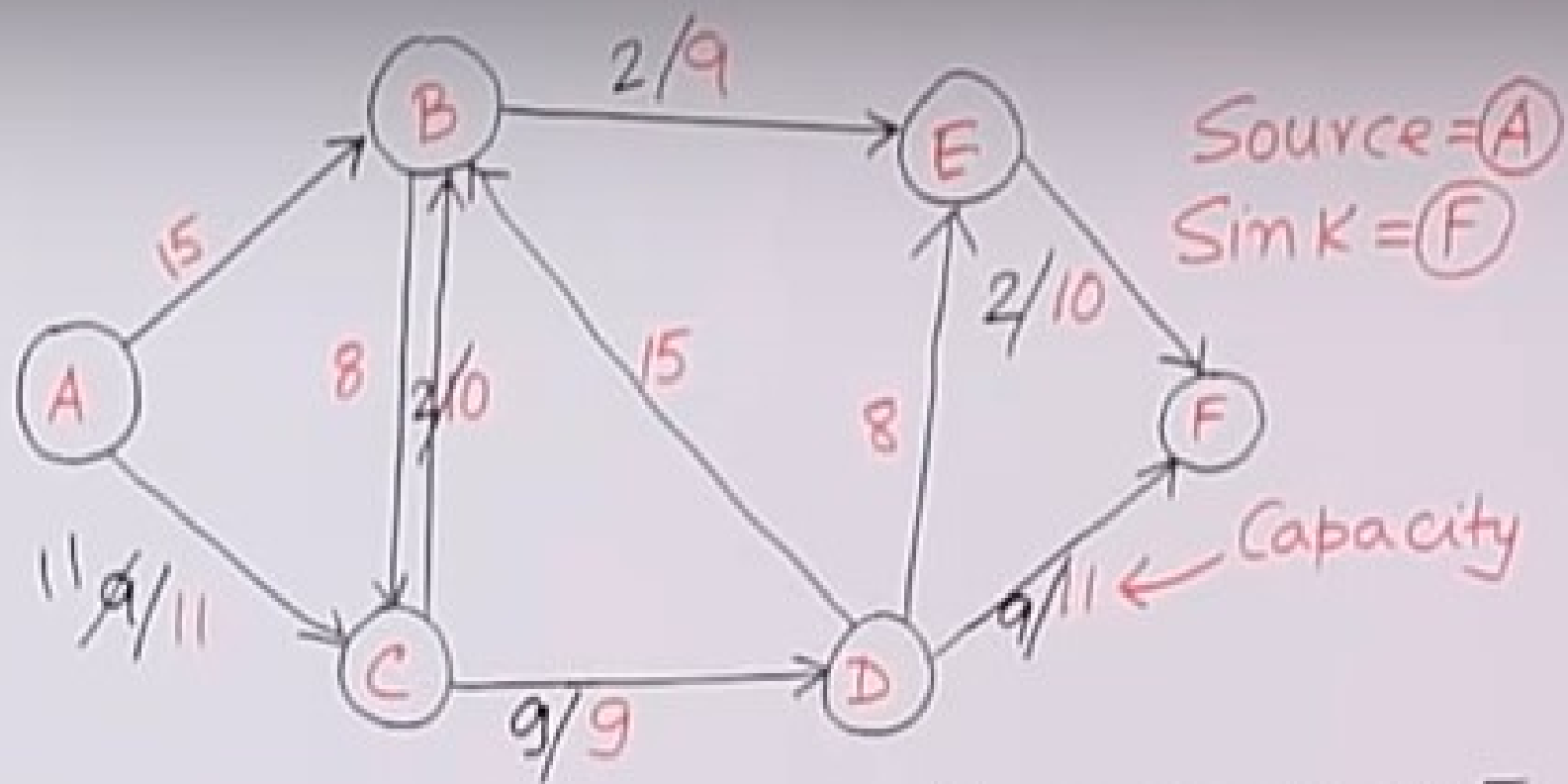
- * **Residual Graph:** It's a graph which indicates additional possible flow. If there is such path from Source to sink then there is a possibility to add flow
- * **Residual Capacity:** It's original capacity of the edge minus flow.
- * **Minimal cut:** Also known as bottle neck capacity, which decides maximum possible flow from Source to Sink through an augmented path
- * **Augmenting path:** Augmenting path can be done in two ways -
 - 1) Non-full forward edges
 - 2) Non-empty backward edges.



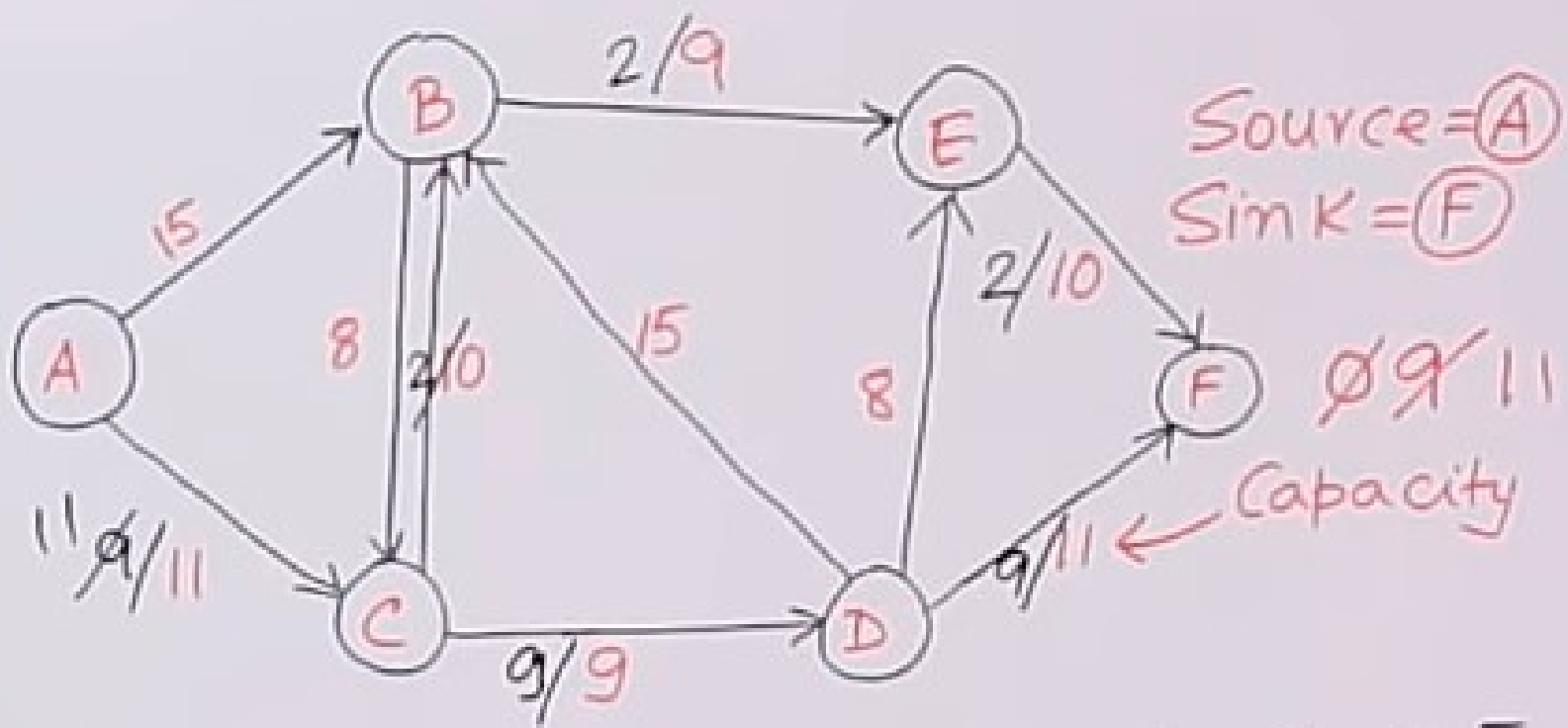




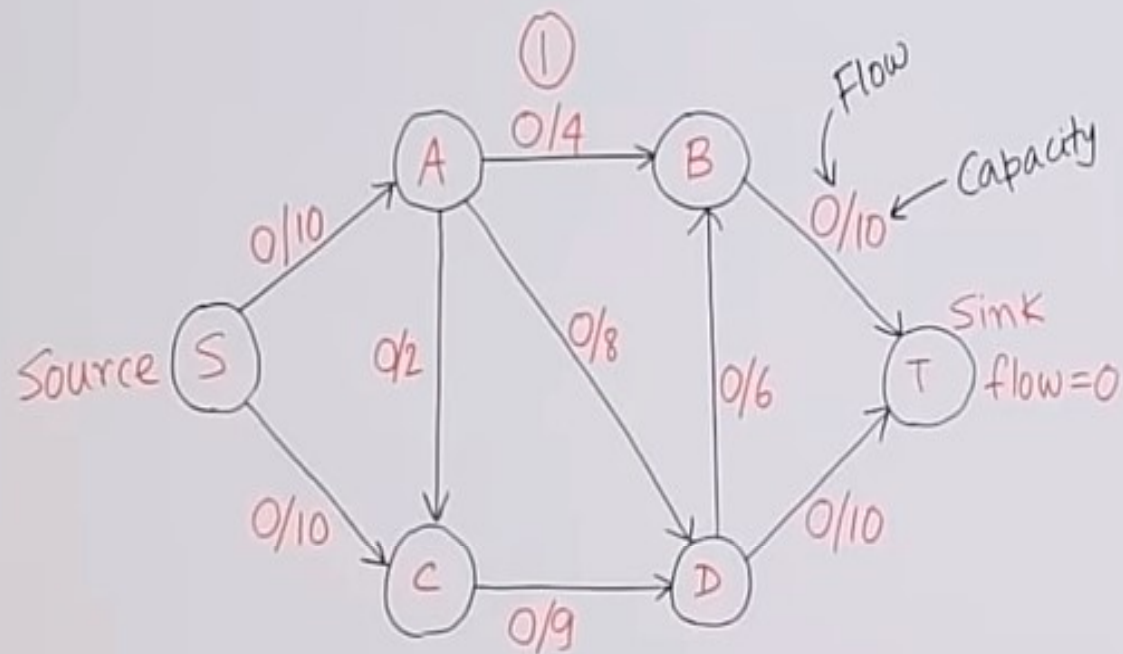
$A \rightarrow C \rightarrow D \rightarrow F$
 $A \rightarrow C \rightarrow B \rightarrow E \rightarrow F$

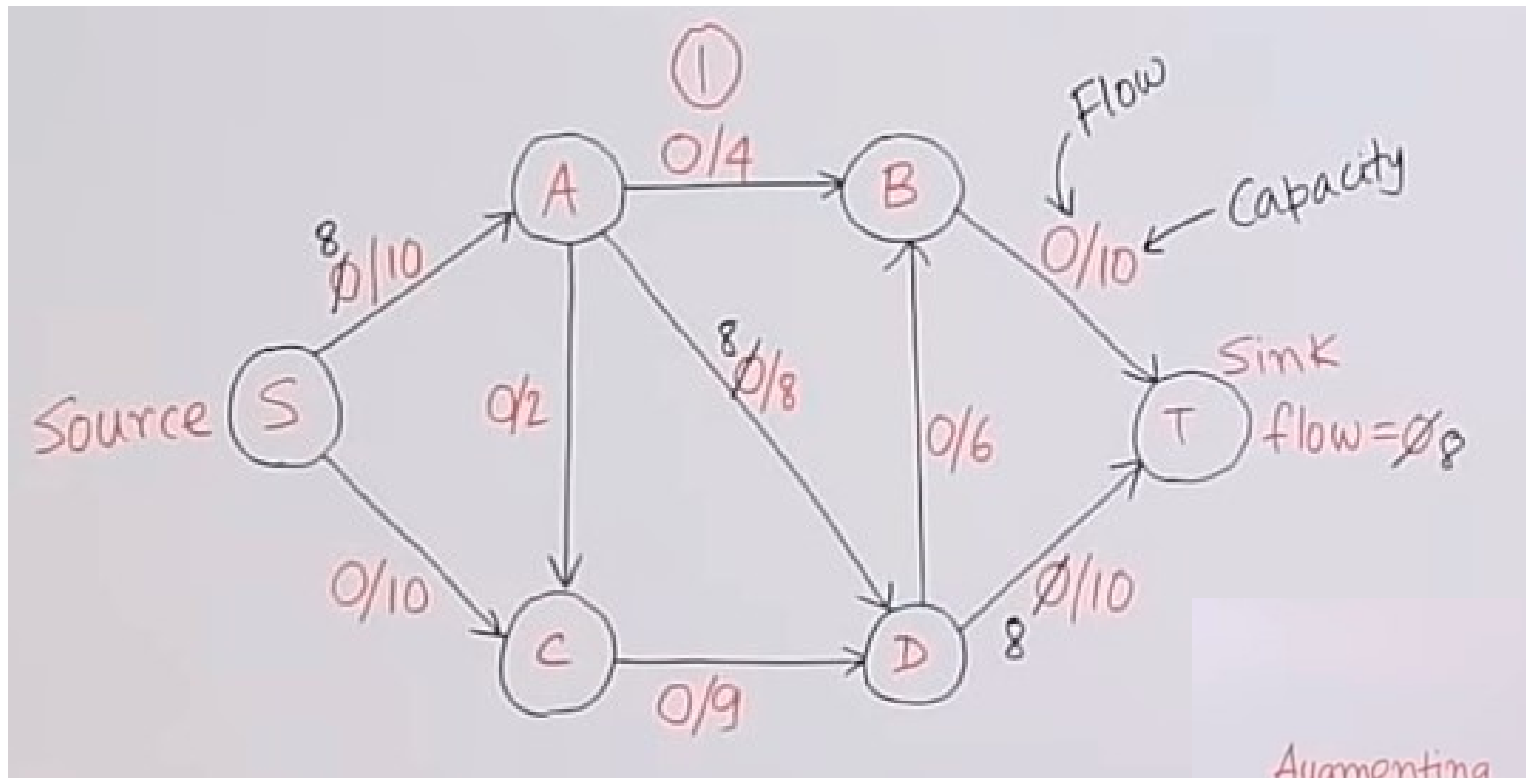


$A \rightarrow C \rightarrow D \rightarrow F$
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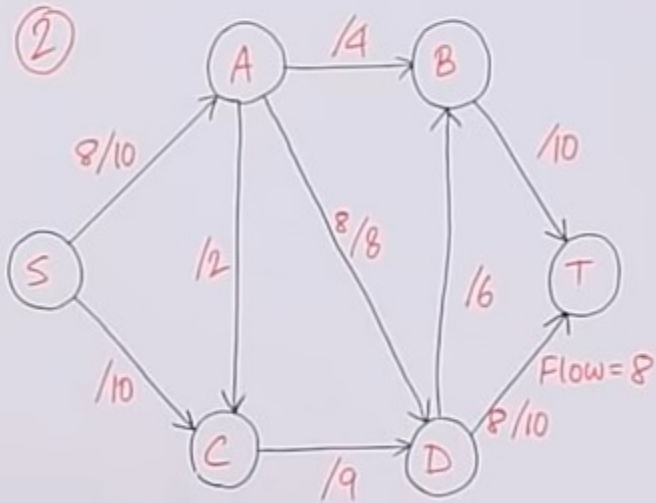
Ford-Fulkerson Algorithm for Maximum Flow Problem: Example





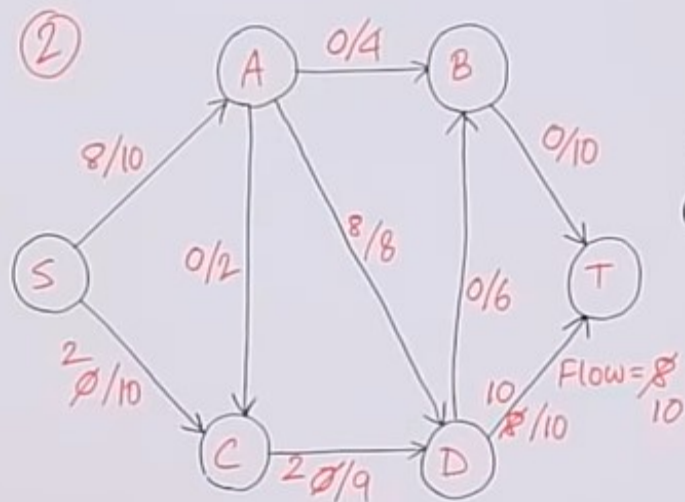
Augmenting paths	Bottle neck capacity
① $S \rightarrow A \rightarrow D \rightarrow T$	8

Problem: Example

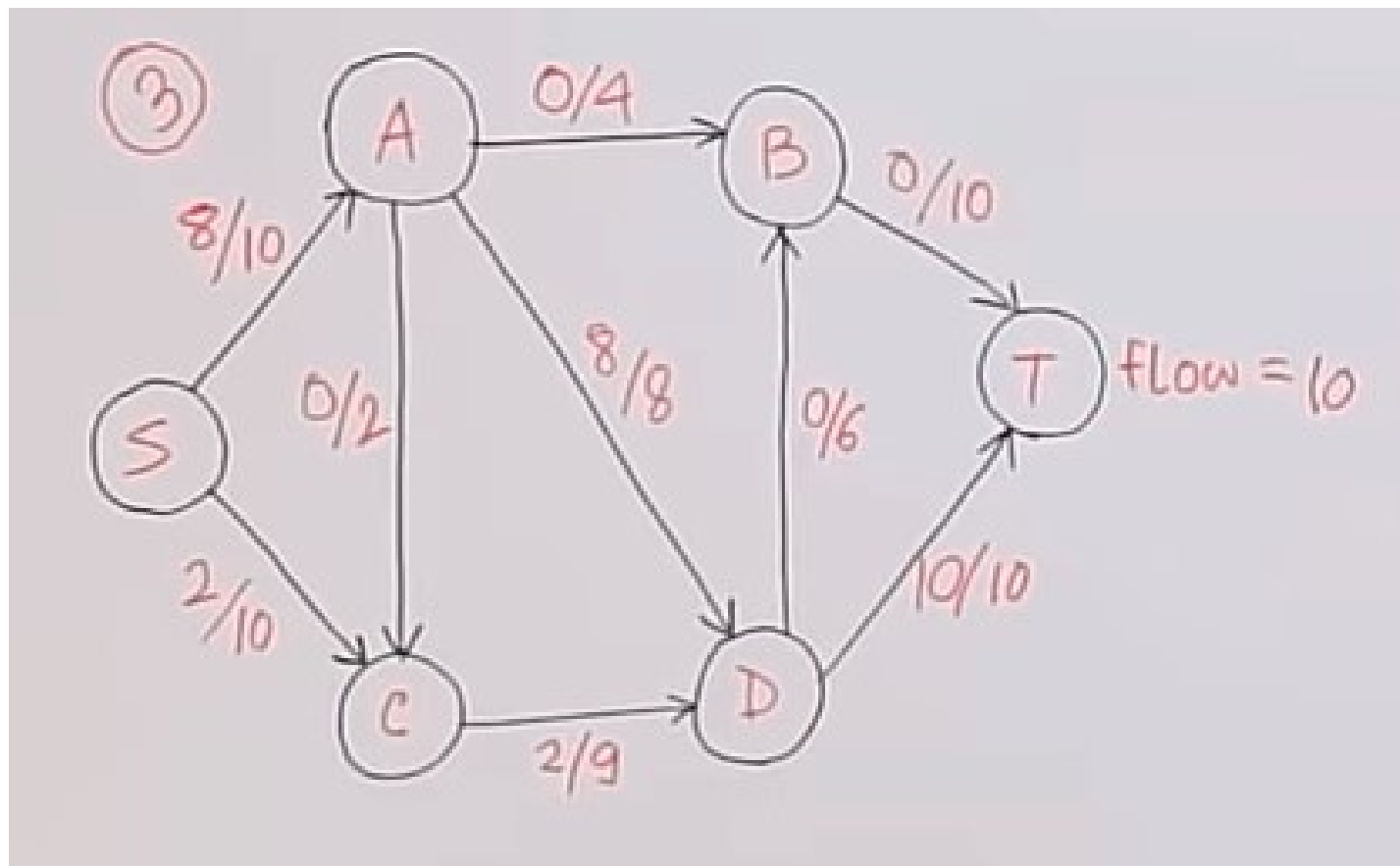


Augmenting paths	Bottle neck capacity
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Problem: Example

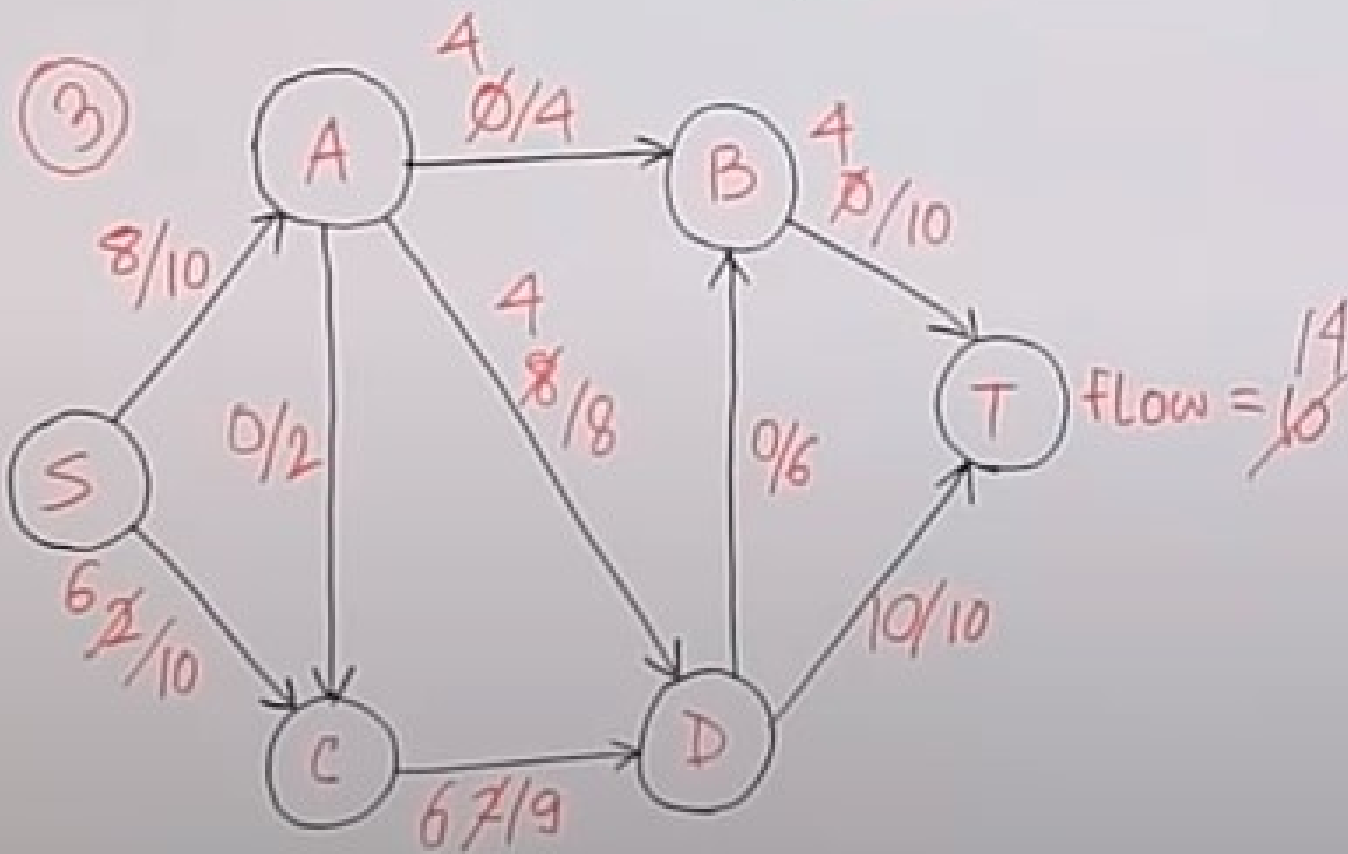


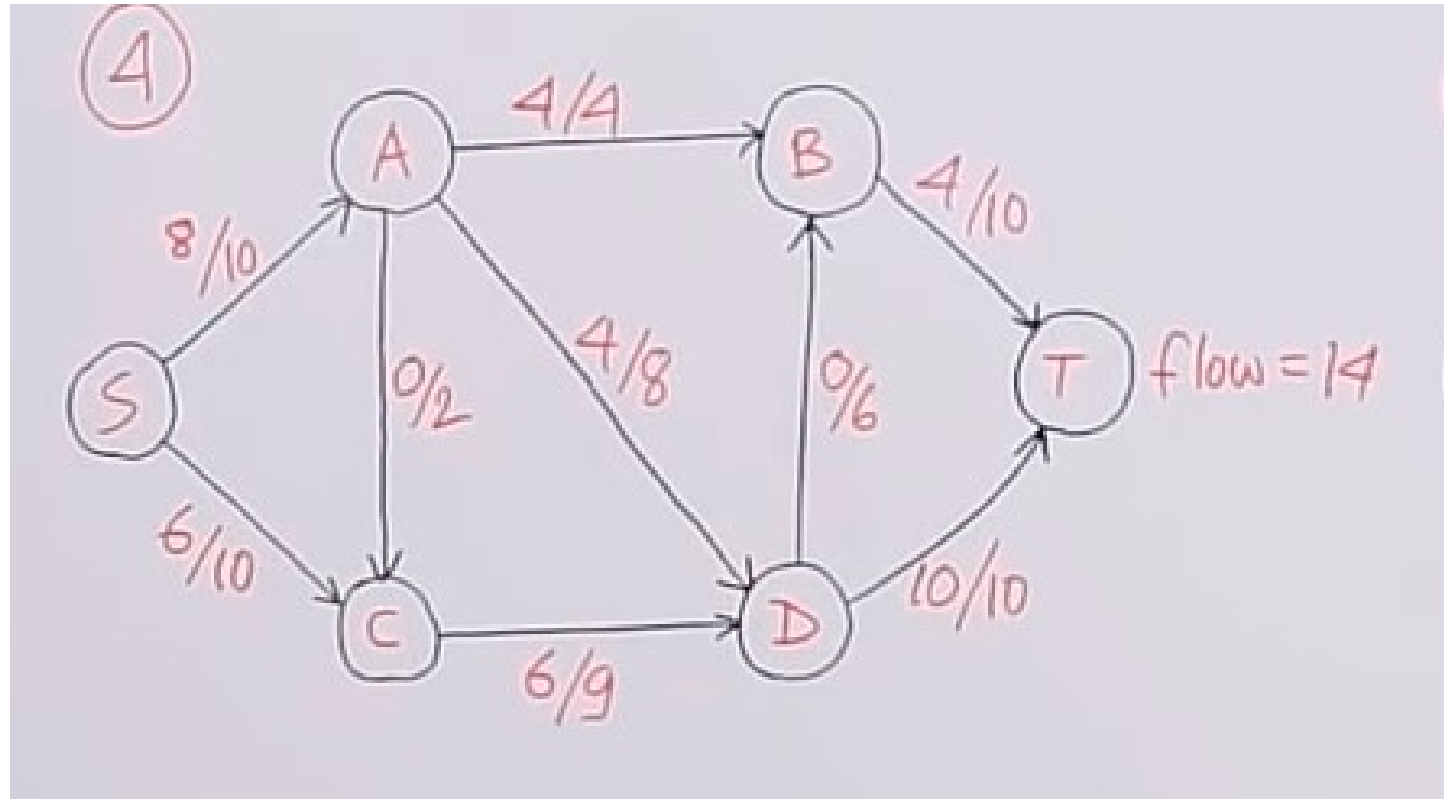
Augmenting paths	Bottle neck capacity
① $S \rightarrow A \rightarrow D \rightarrow T$	8
② $S \rightarrow C \rightarrow D \rightarrow T$	2



Augmenting paths	Bottle neck capacity
① $S \rightarrow A \rightarrow D \rightarrow T$	8
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③ $S \rightarrow C \rightarrow D \rightarrow A \rightarrow B \rightarrow T$	

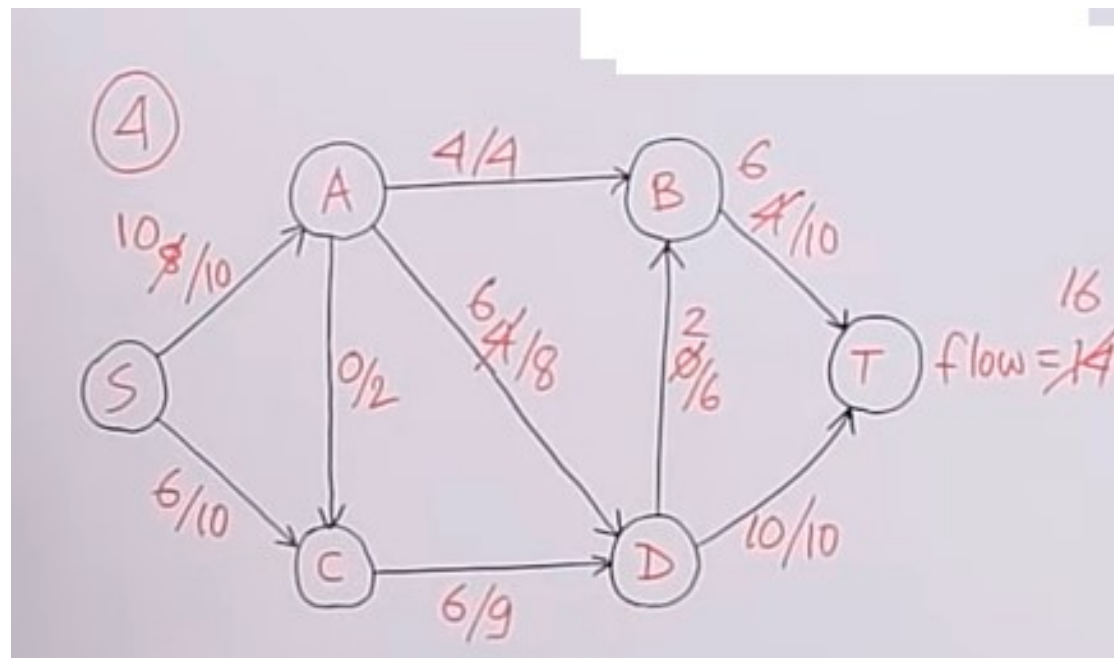
Augmenting paths	Bottle neck capacity
① $S \rightarrow A \rightarrow D \rightarrow T$	8
② $S \rightarrow C \rightarrow D \rightarrow T$	2
③ $S \rightarrow C \rightarrow D \rightarrow A \rightarrow B \rightarrow T$	4

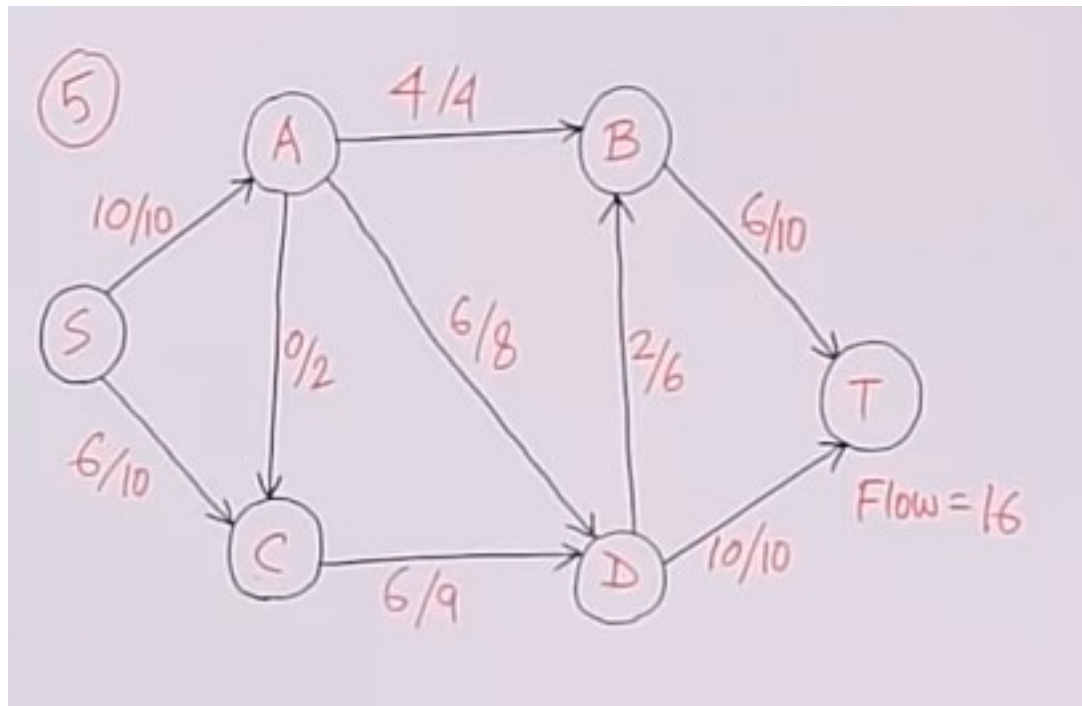




Augmenting paths	Bottle neck capacity
① $S \rightarrow A \rightarrow D \rightarrow T$	8
② $S \rightarrow C \rightarrow D \rightarrow T$	2
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④ $S \rightarrow A \rightarrow D \rightarrow B \rightarrow T$	2

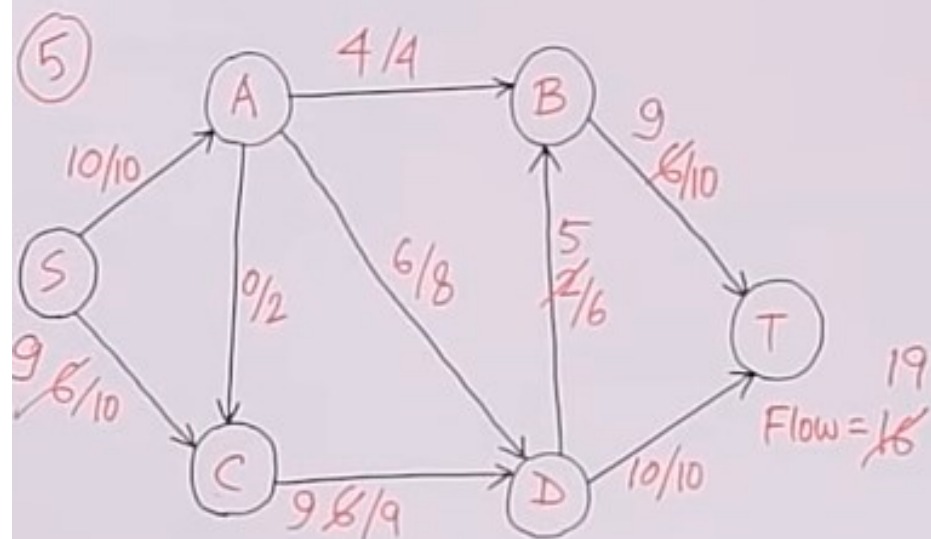
$\frac{8}{10}$

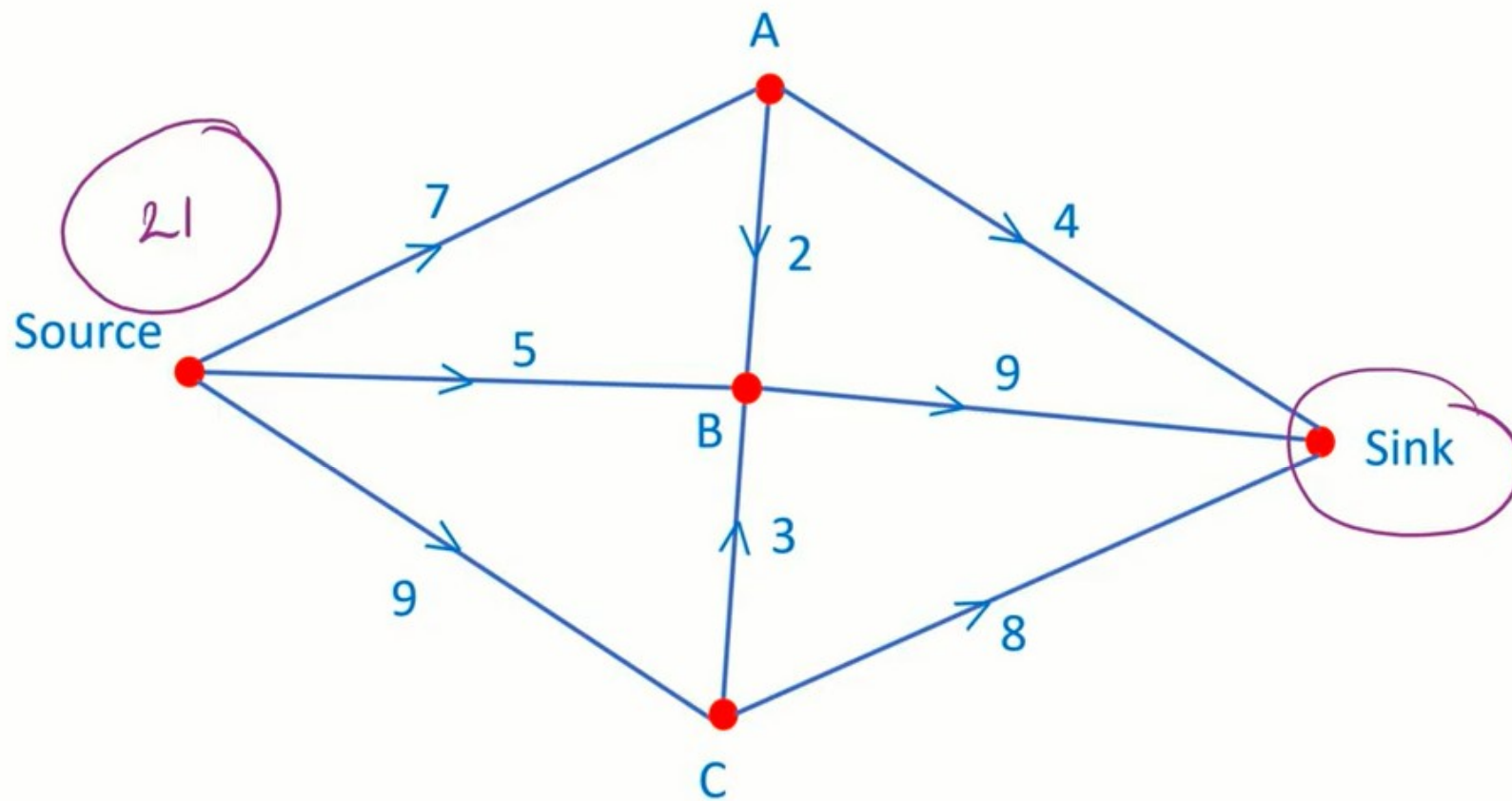


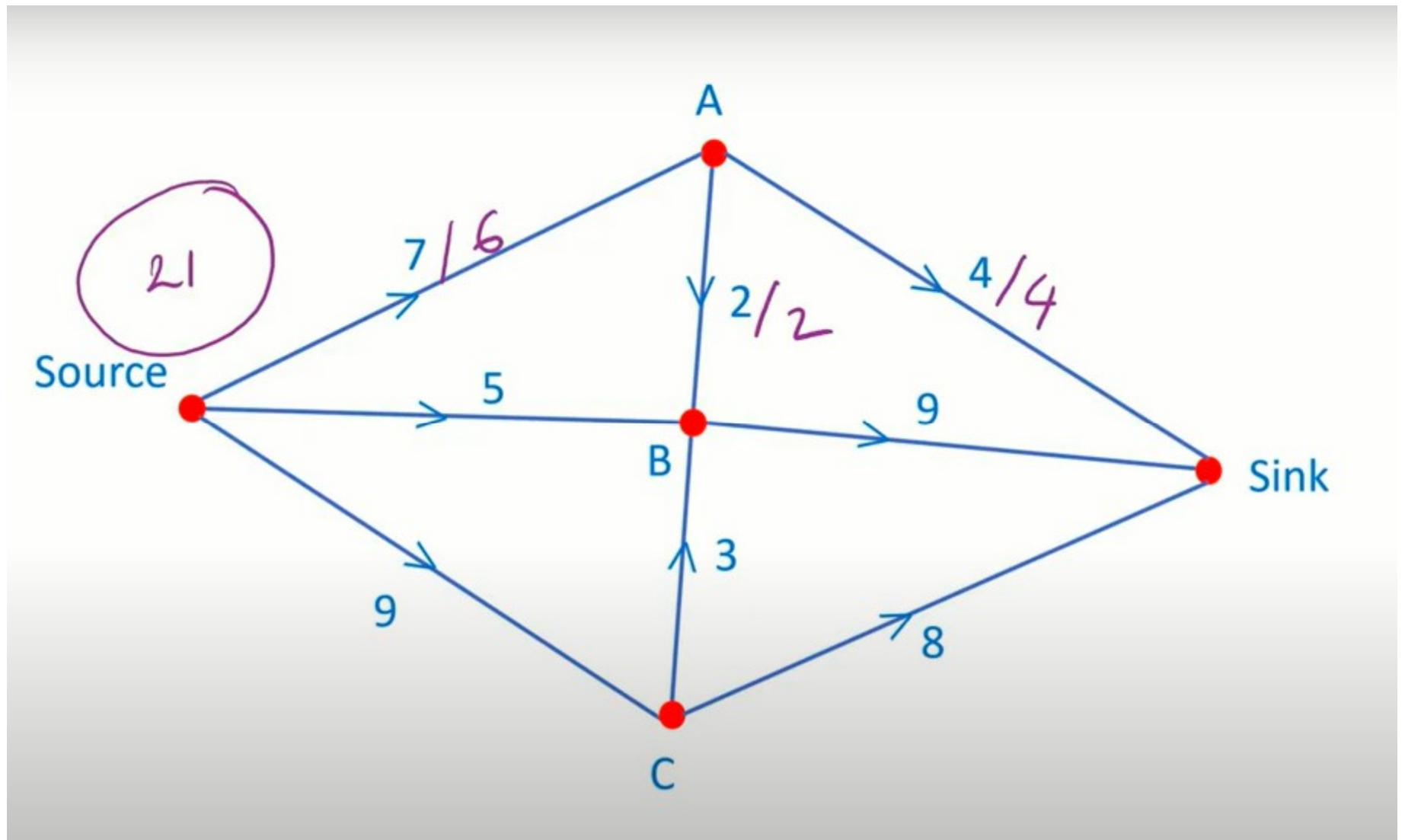


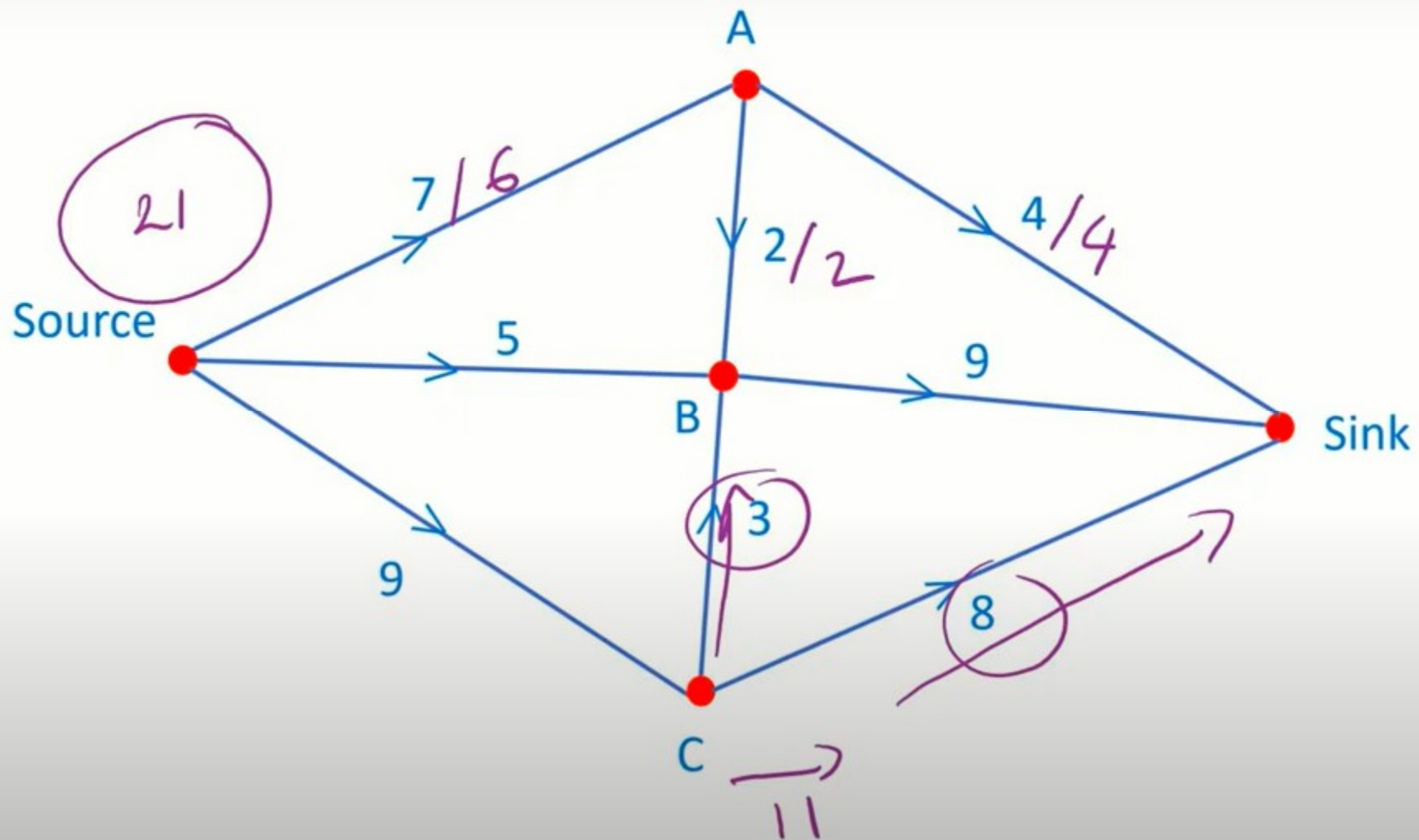
Augmenting paths	Bottle neck capacity
① $S \rightarrow A \rightarrow D \rightarrow T$	8
② $S \rightarrow C \rightarrow D \rightarrow T$	2
③ $S \rightarrow C \rightarrow D \rightarrow A \rightarrow B \rightarrow T$	4
④ $S \rightarrow A \rightarrow D \rightarrow B \rightarrow T$	2
④ ⑤ $S \rightarrow C \rightarrow D \rightarrow B \rightarrow T$	3

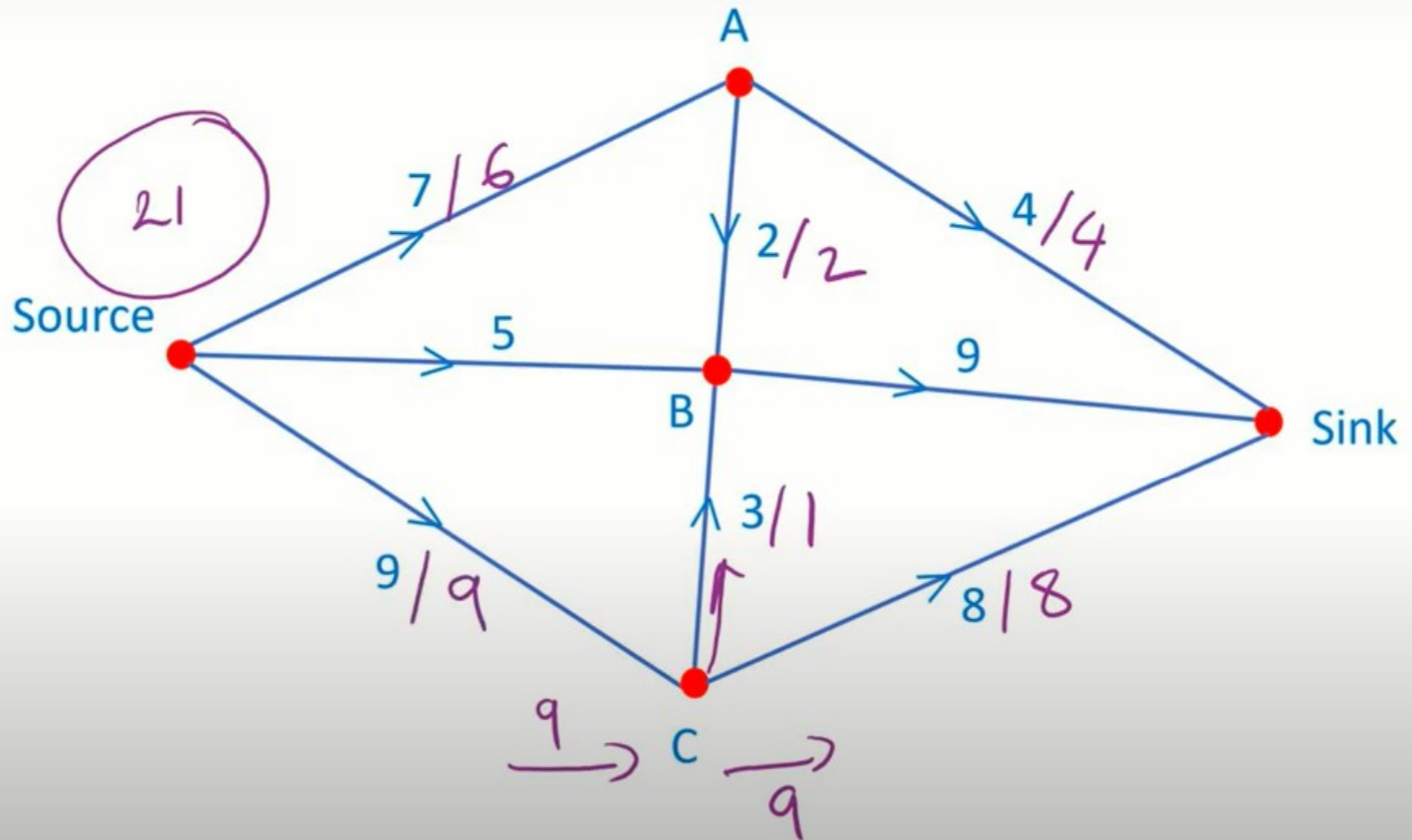
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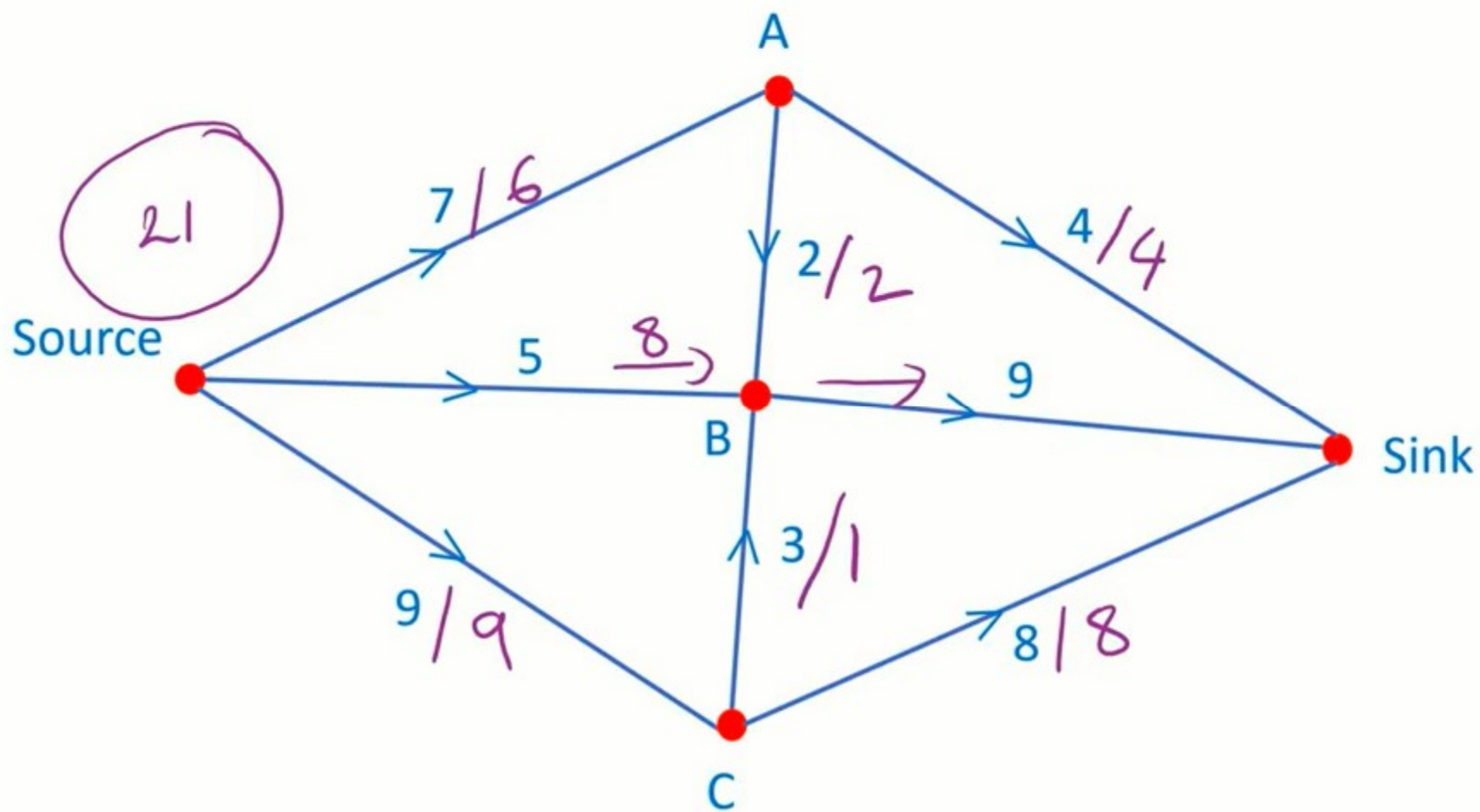


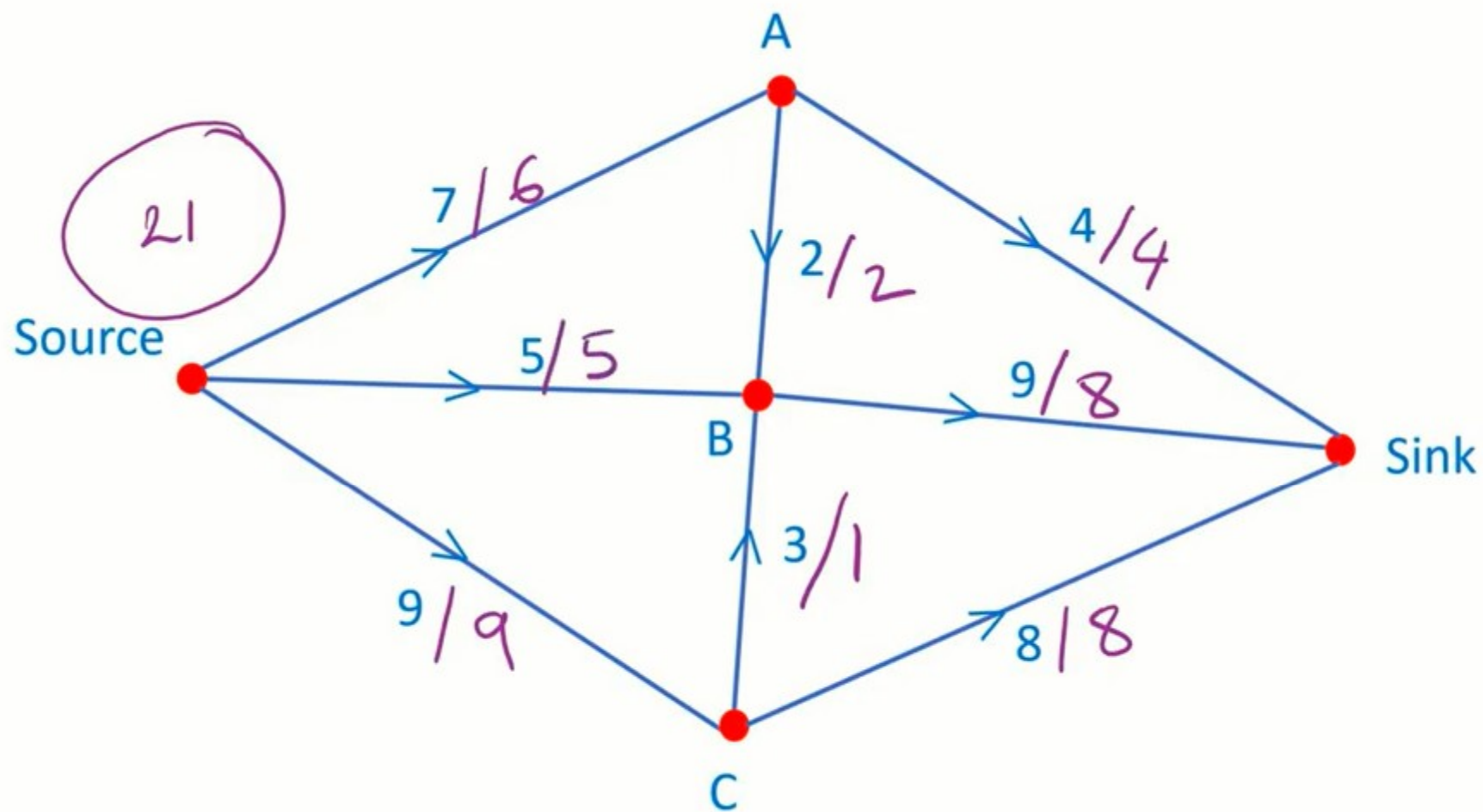


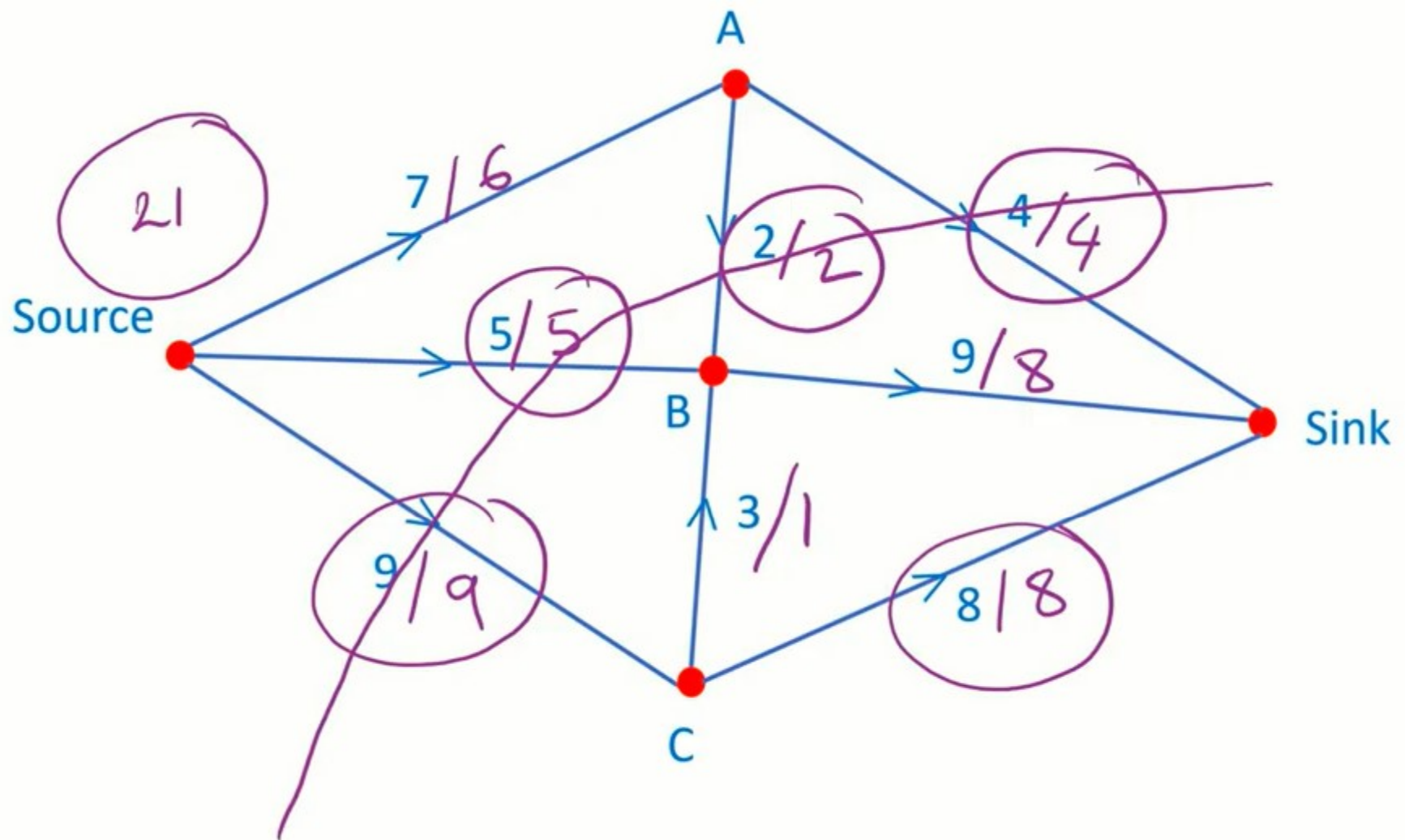




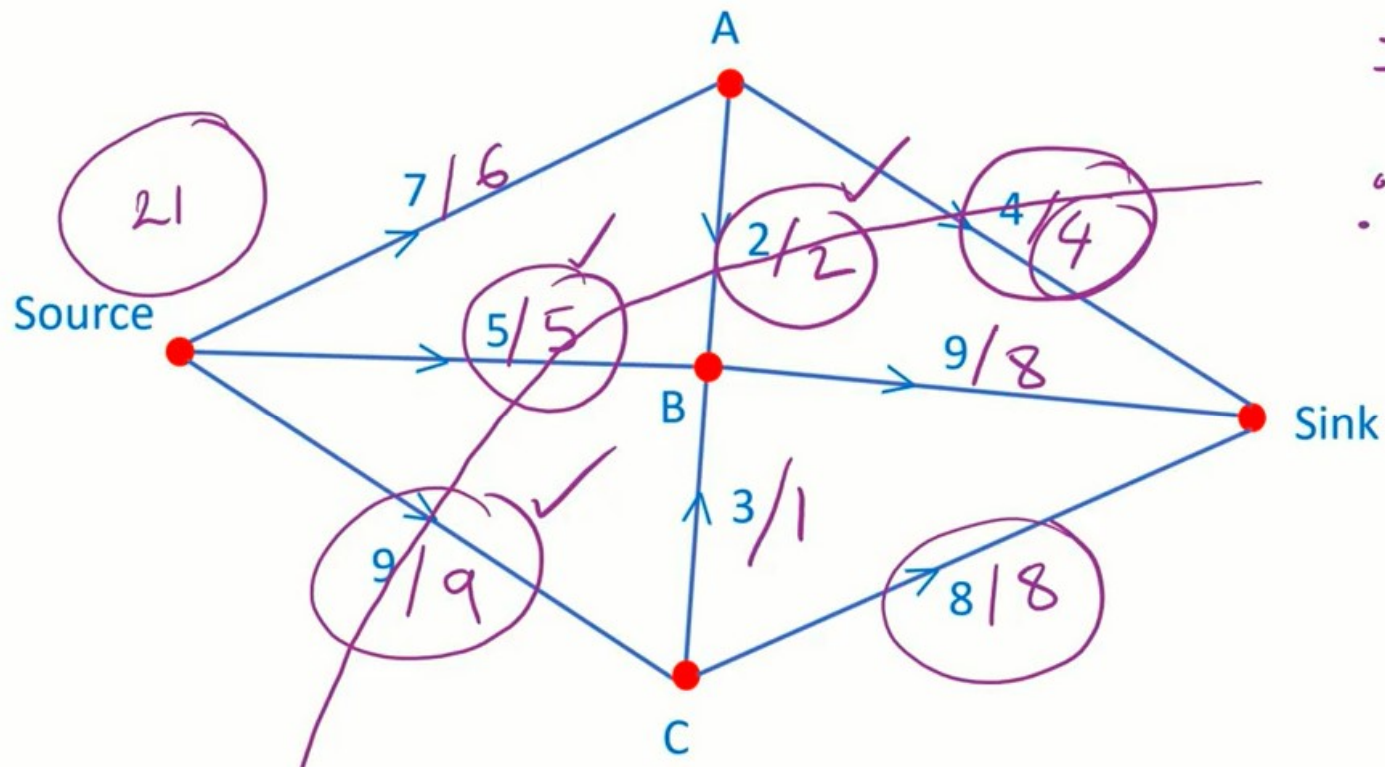








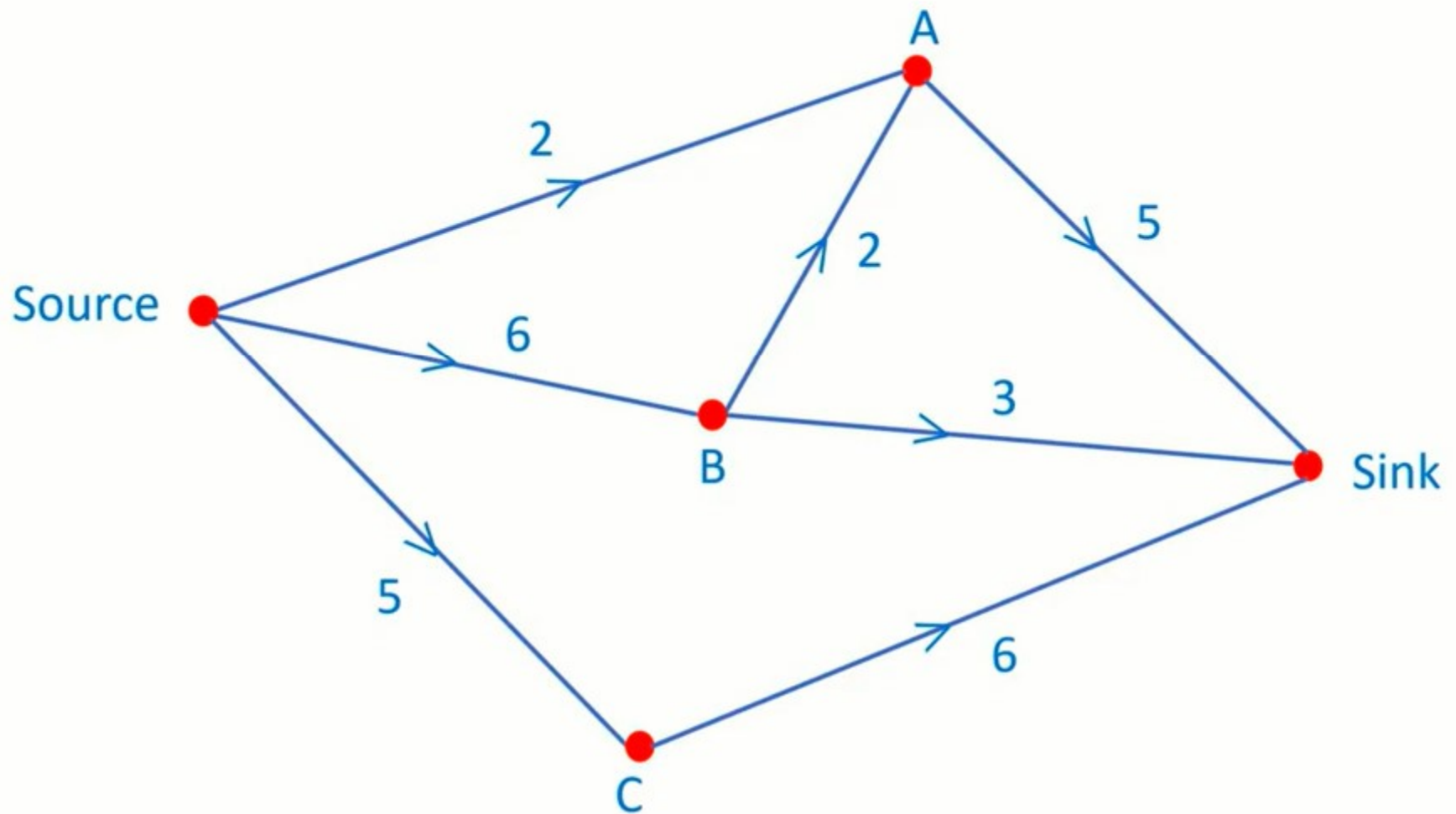
Find the maximum flow for the network diagram below.



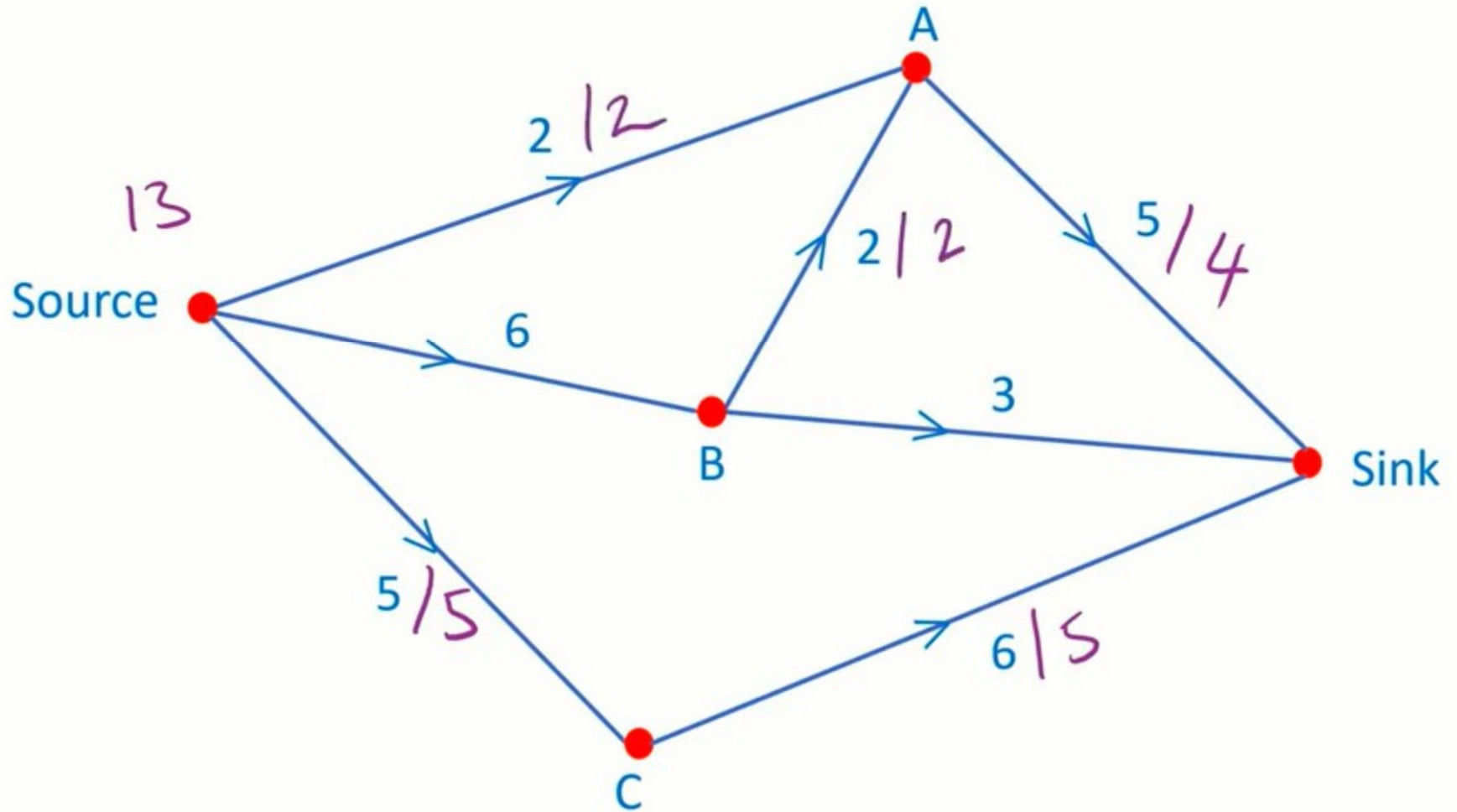
$$4 + 2 + 5 + 9 = 20$$

$\therefore$ the maximum flow is 20

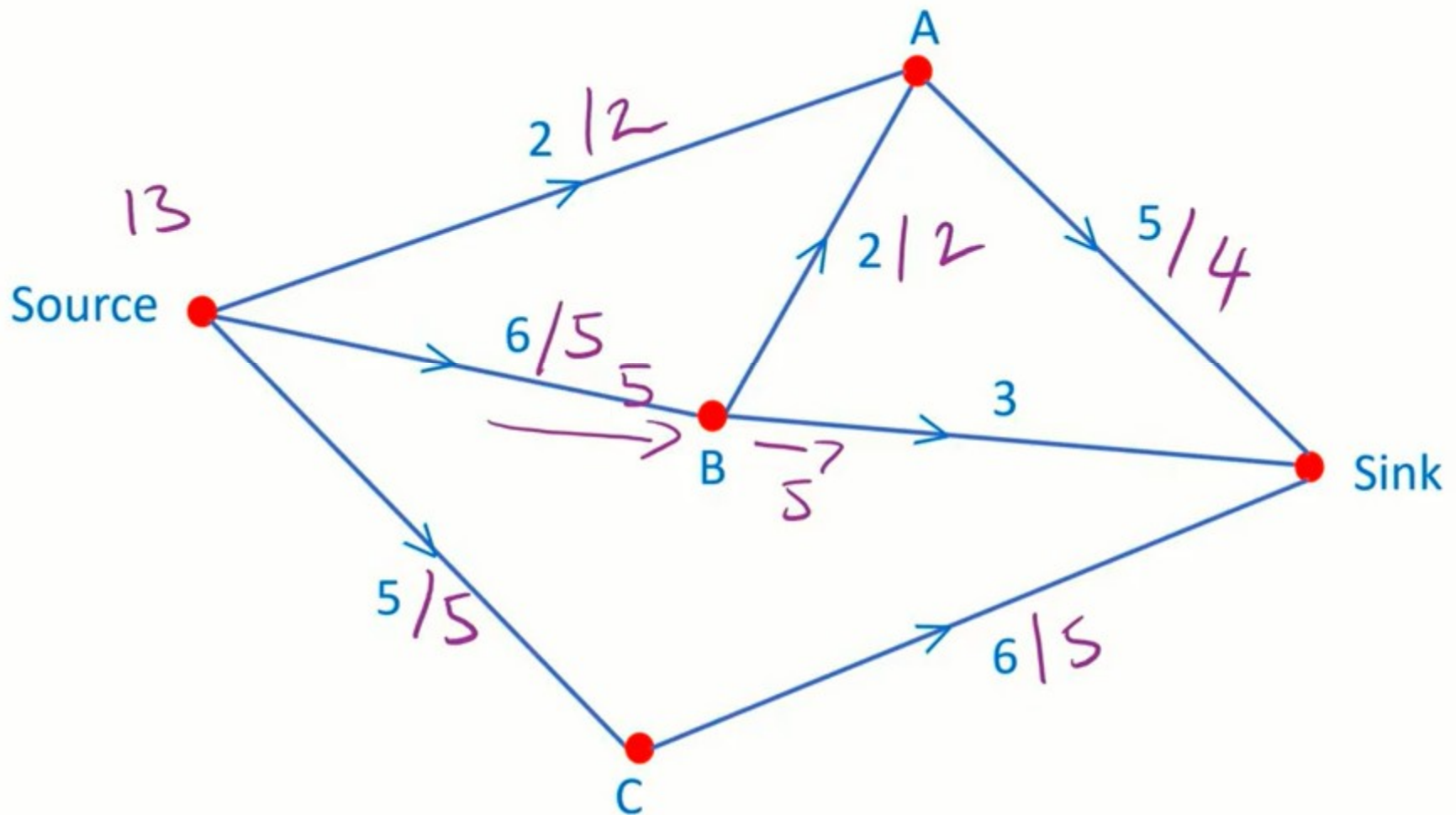
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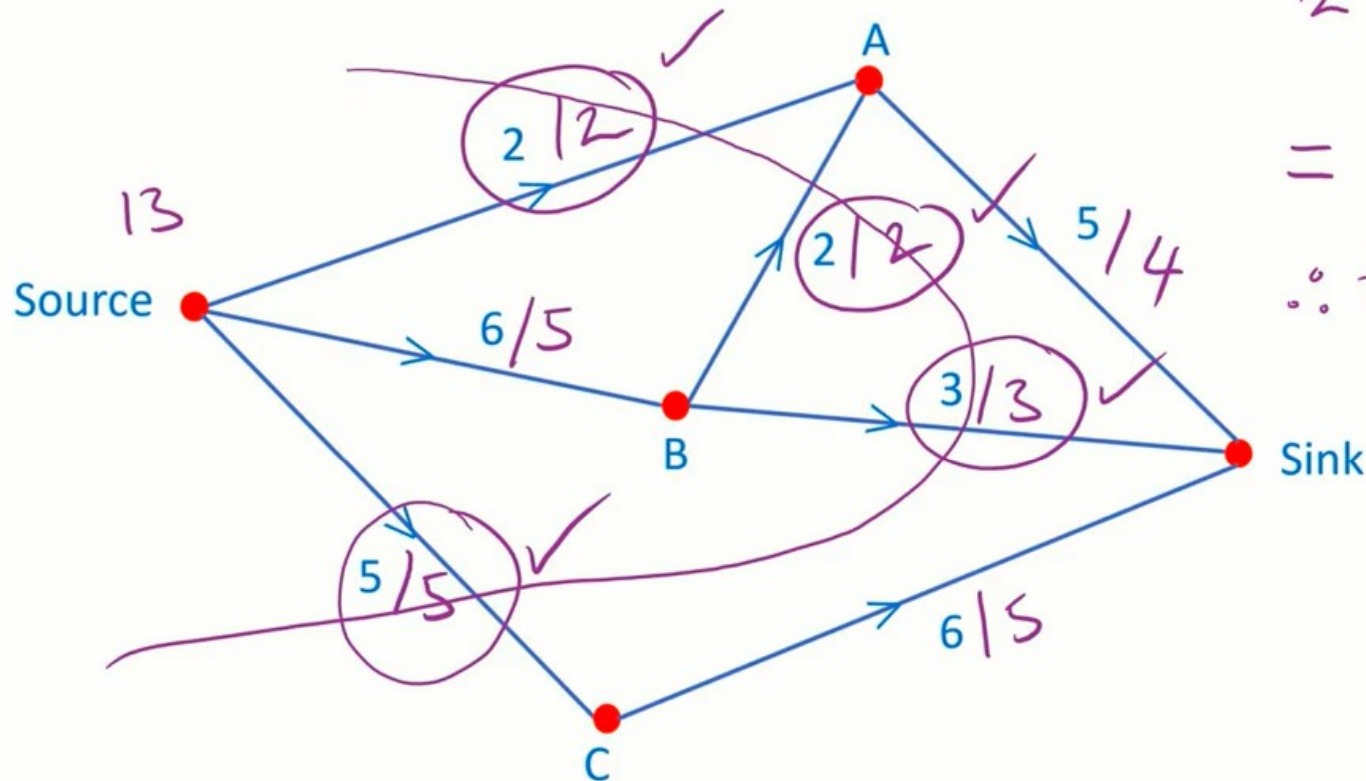
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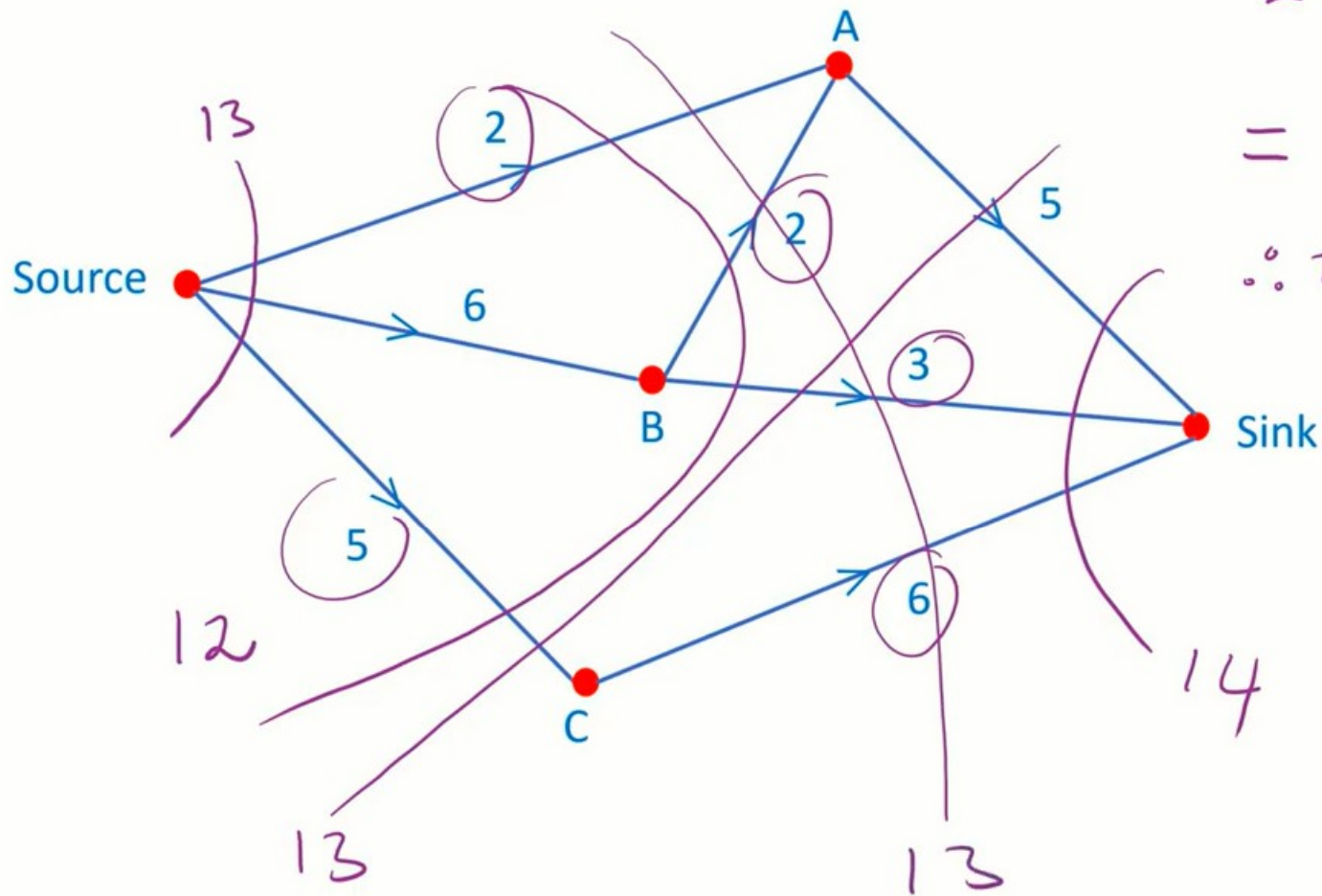
Find the maximum flow for the network diagram below.



$$2 + 2 + 3 + 5 = 12$$

$\therefore$ the maximum flow is 12

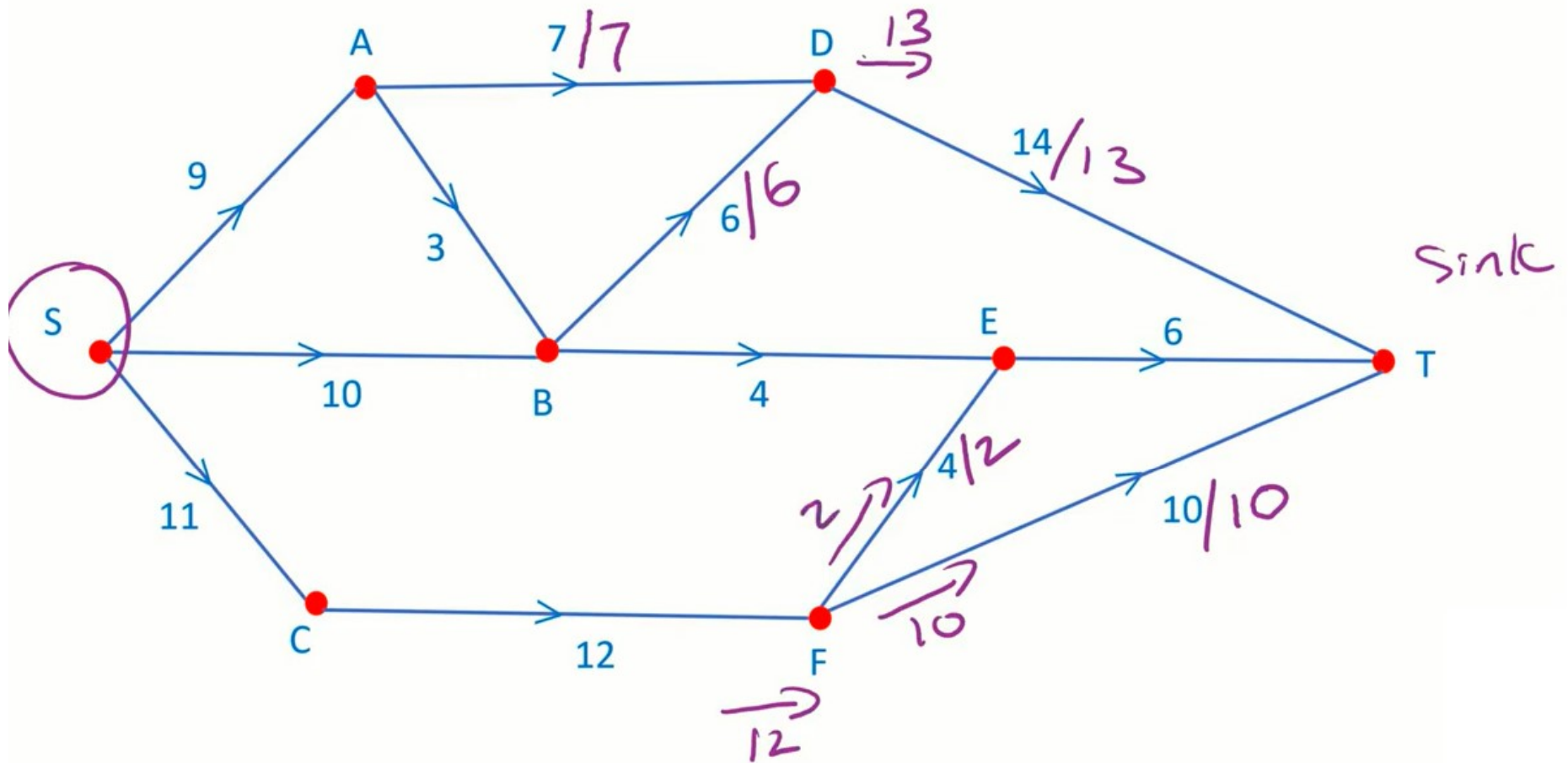
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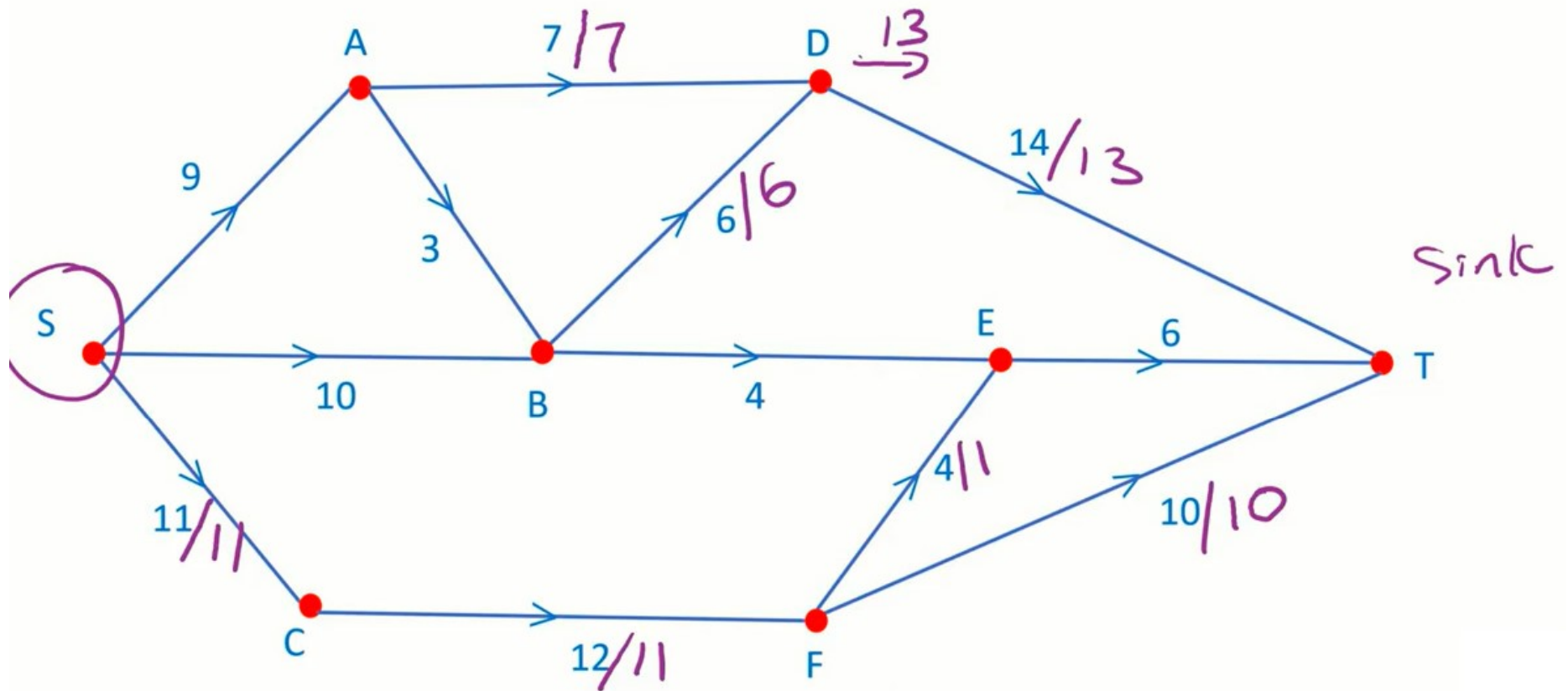
$$2 + 2 + 3 + 5 = 12 \checkmark \checkmark$$

$\therefore$ the maximum flow is 12

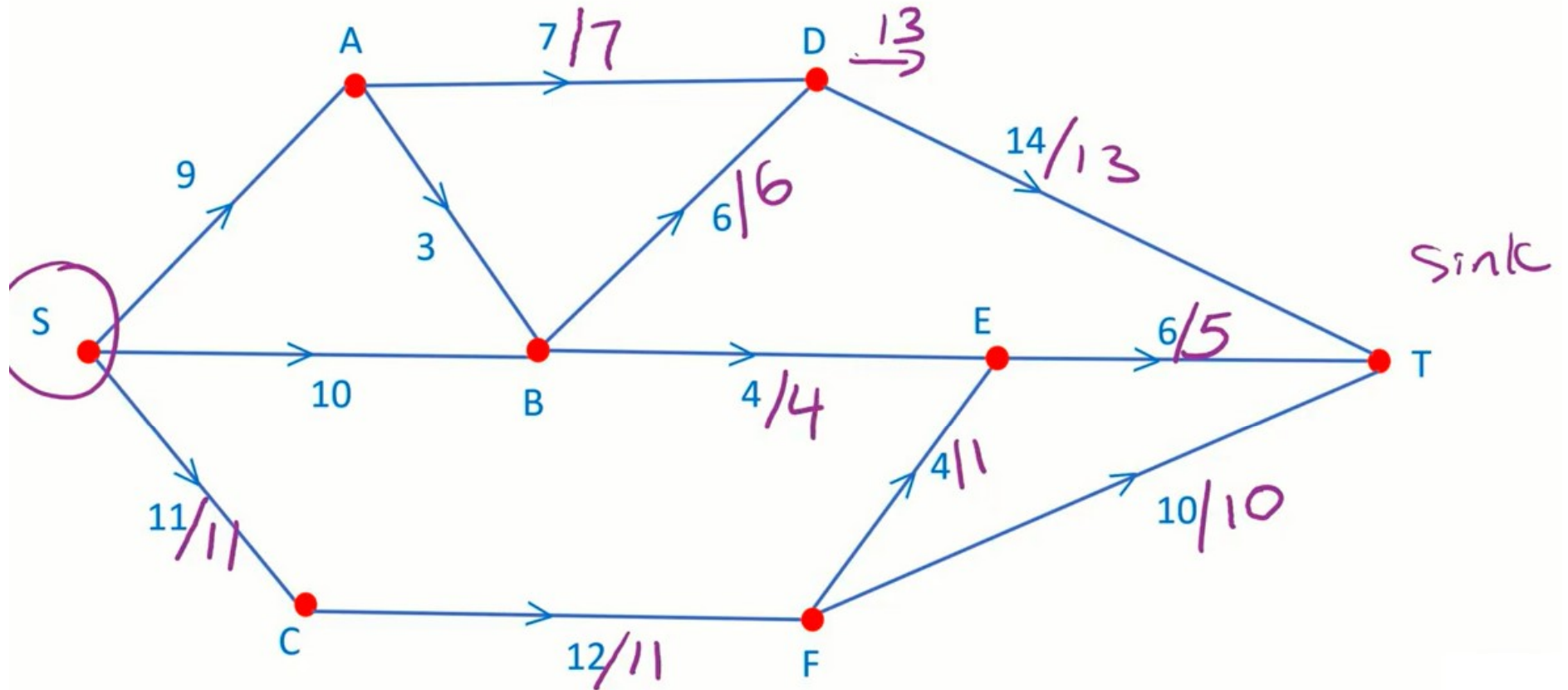
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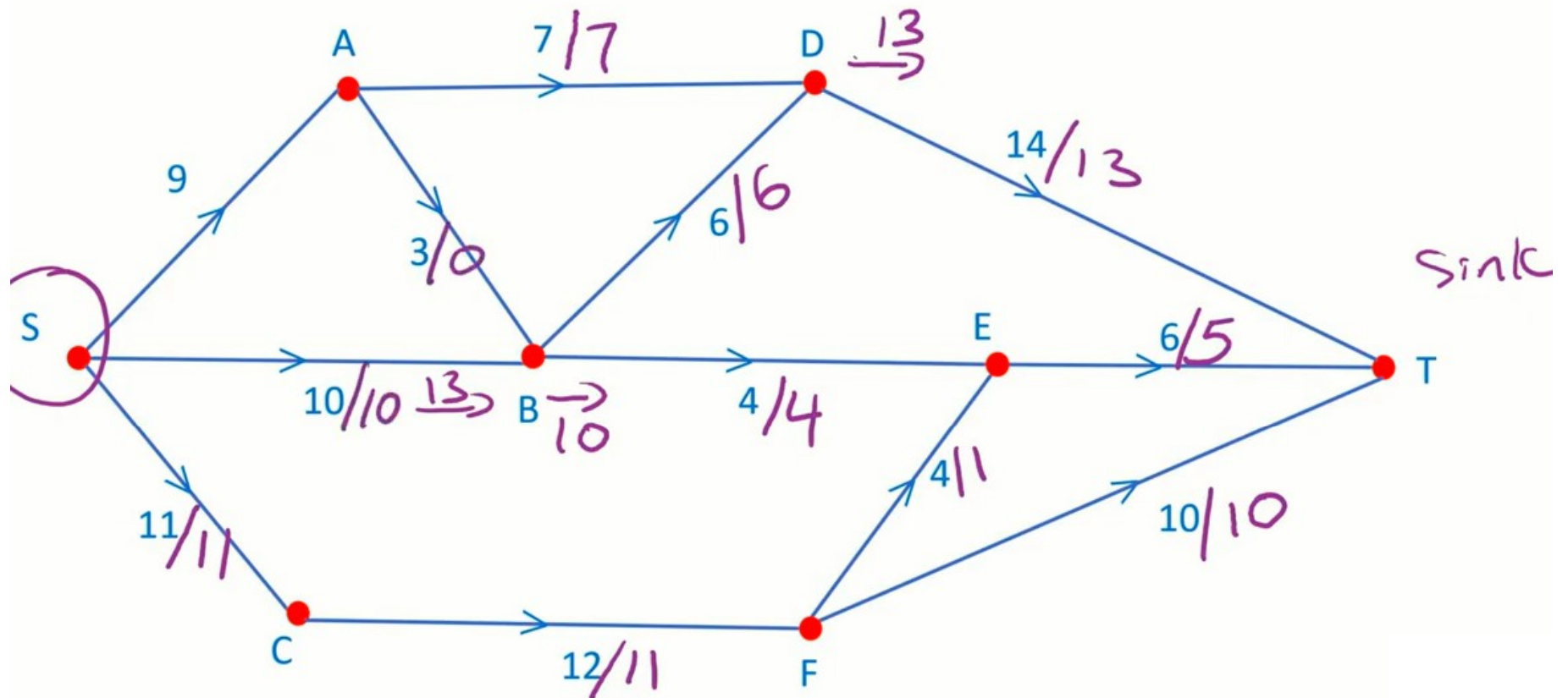
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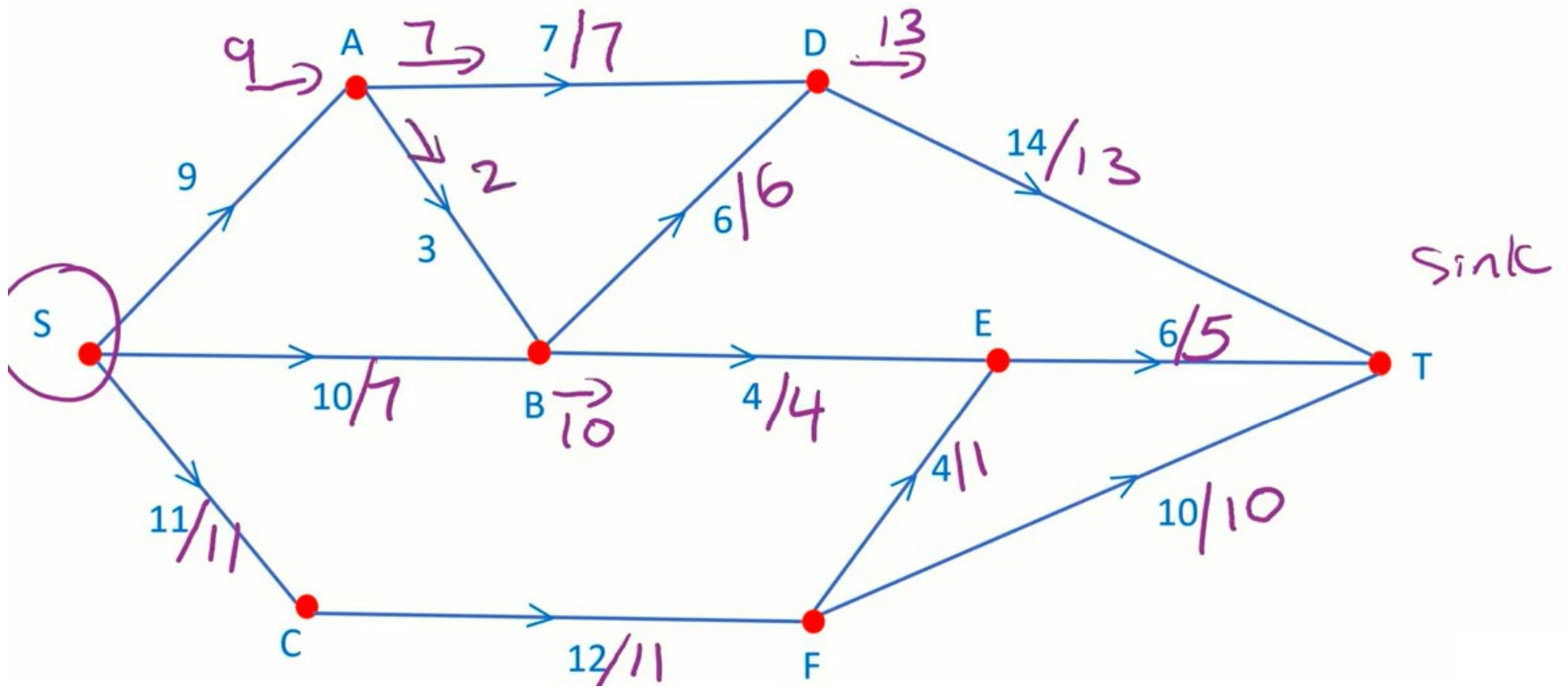
Find the maximum flow for the network diagram below.



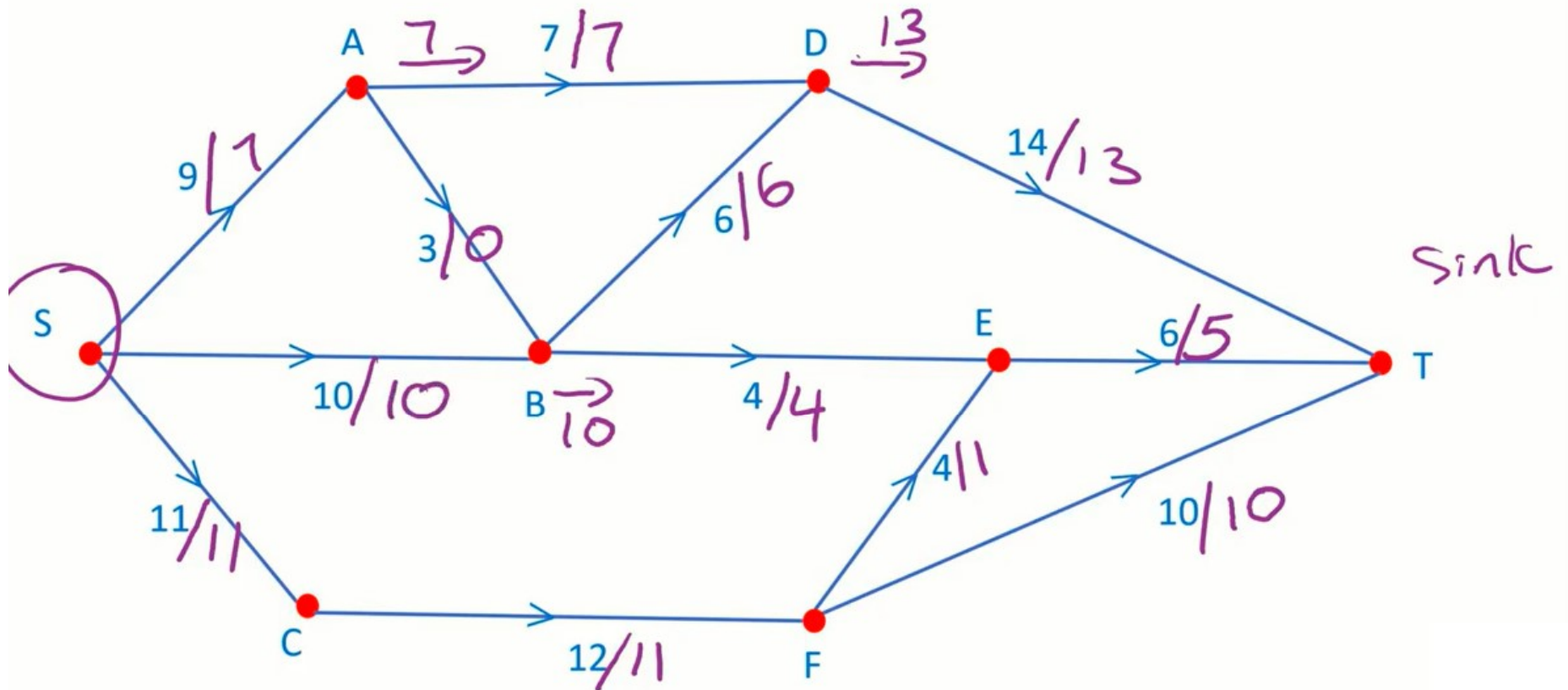
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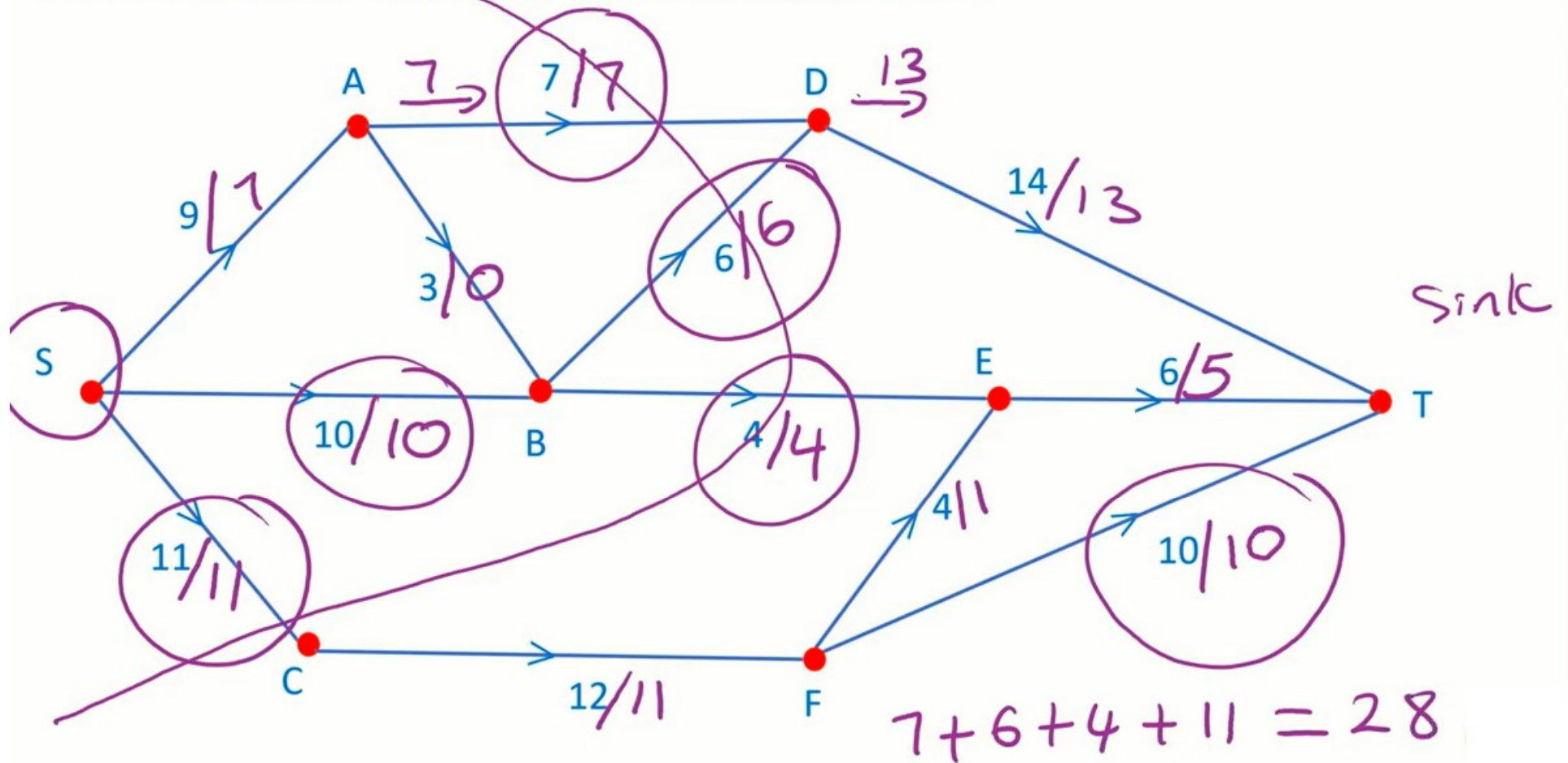
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